

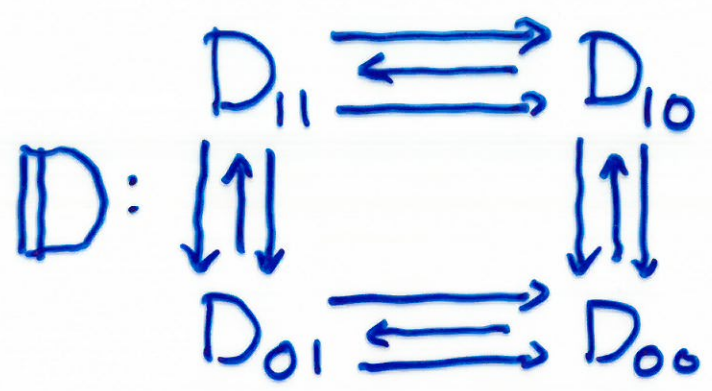
Free Double Categories and the Word Problem for Groups

by

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Toronto, September 24, 2000

WANT TO DESCRIBE THE FREE
 DOUBLE CATEGORY GENERATED
 BY A DOUBLE REFLEXIVE GRAPH

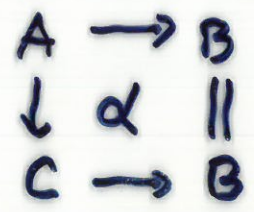


- DOUBLE GRAPH, CAT GRAPH → TOO EASY
- REFLEXIVE GRAPH IN CAT → NOT PRACTICAL

• LEFT ADJ EXISTS $DCAT \xrightleftharpoons[F]{U} DGPH$

DESCRIBE **EXPLICITLY**

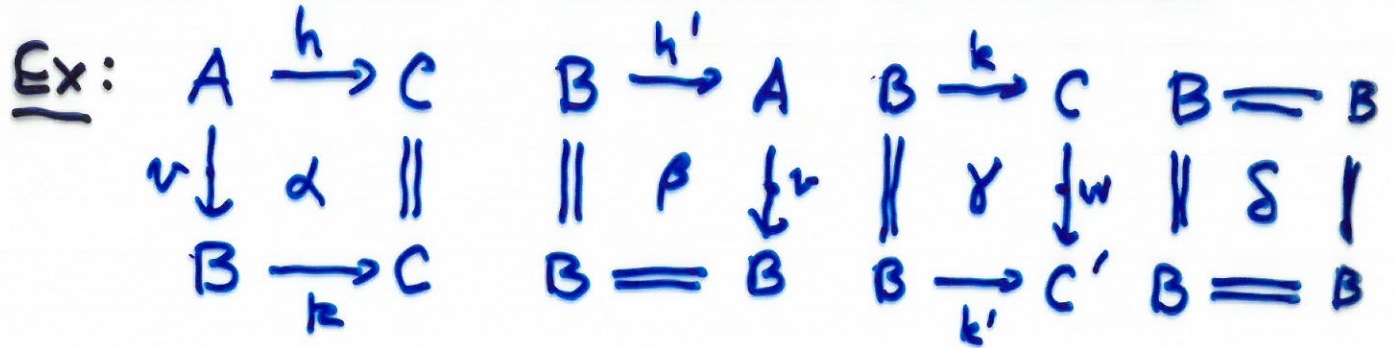
• TYPICAL CELL IN ID



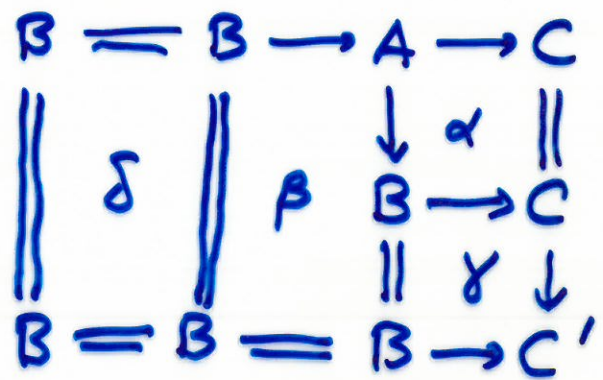
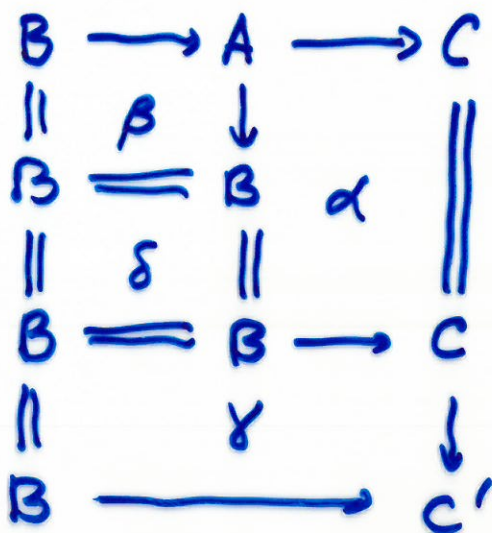
IDENTITIES ON BOUNDARY
 ⇒ EQUATIONS IN FREE!

F(D) : OBJ, HOR ARR, VERT ARROW - EASY

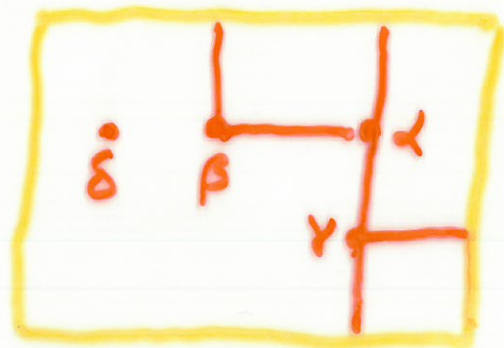
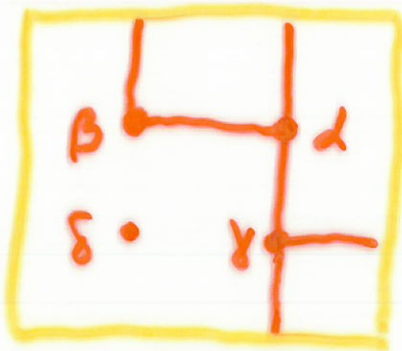
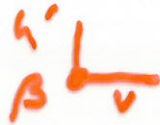
CELLS = EQUIV CLASSES OF WORDS



$$((\beta \circ_v \delta) \circ_h \alpha) \circ_v \gamma = (\delta \circ_h \beta) \circ_h (\alpha \circ_v \gamma)$$



PLANAR
DUAL'S



"ECKMANN - HILTON"

RECTANGULAR COMPLEXES

- FIRST STUDY THE SHAPE OF CELLS

$$K = (R, P, H, V)$$

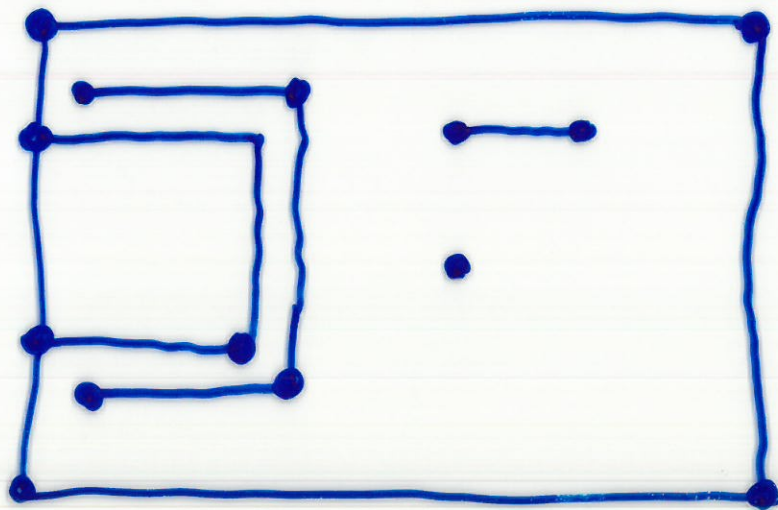
$$R = [a, b] \times [c, d] \subseteq \mathbb{R}^2$$

P = FIN SET OF POINTS — VERTICES

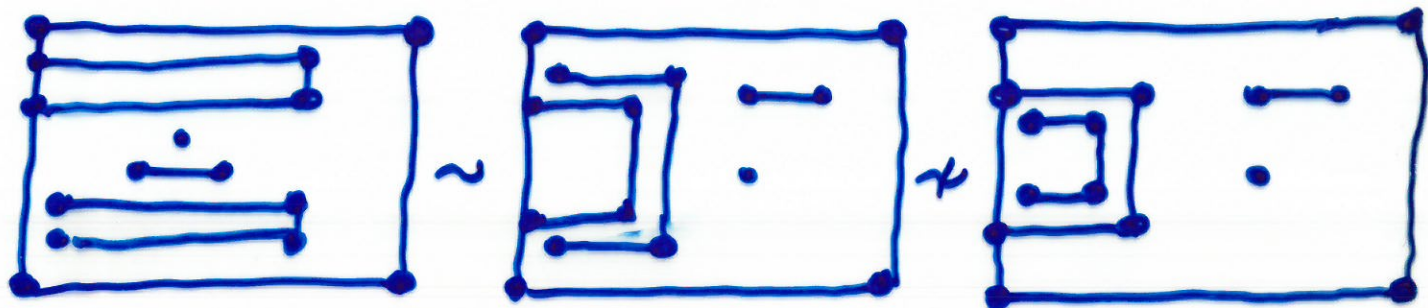
H = FIN SET OF HOR SEGMENTS

V = " " " VERT " " } EDGES

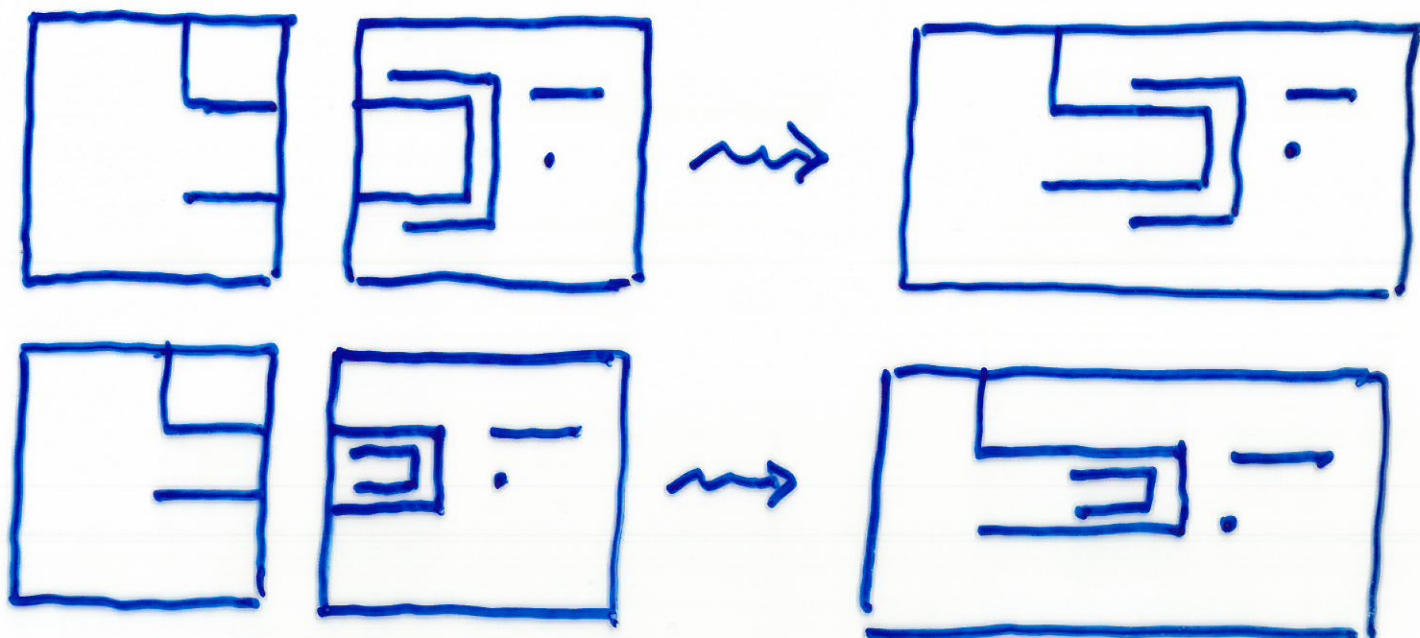
1. END PTS OF EDGES ARE VERTICES
2. EVERY PT OF ∂R IS ON AN EDGE
3. EDGES ^{INTERSECT} ONLY IN END POINTS
4. ANY VERTEX ON ∂R ENP PT OF INTERIOR ^{EDGE}



TWO COMPLEXES ARE **HOMOTOPIC** IF
 EACH CAN BE DEFORMED CONTINUOUSLY ^{INTO THE OTHER}
 IN A WAY RESPECTING 1-4.

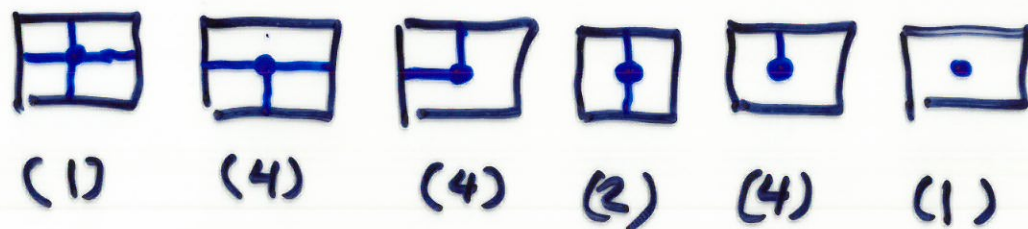


HORIZONTAL MERGING

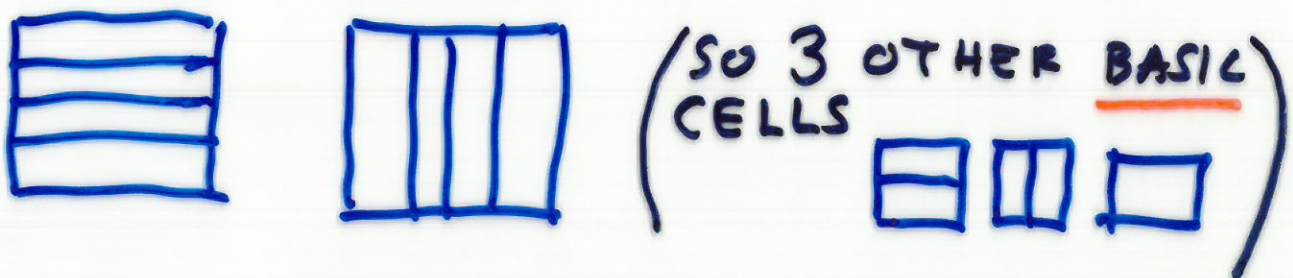


- COMMON BOUNDARY ERASED
- RESPECTS HOMOTOPY

THM 1: HOMOTOPY CLASSES OF RECT COMPLEXES ARE THE CELLS OF A DOUB CAT FREE ON THE DOUB REFL GPH WITH 1 VERTEX, 1 NON IDENTITY HOR (VERT) ARROW AND 16 NON IDENTITY CELLS ONE FOR EACH CHOICE OF IDENTITY OR NON IDENTITY FOR THE ARROWS OF THE BOUNDARY. THEY ARE REPRESENTED BY THE COMPLEXES



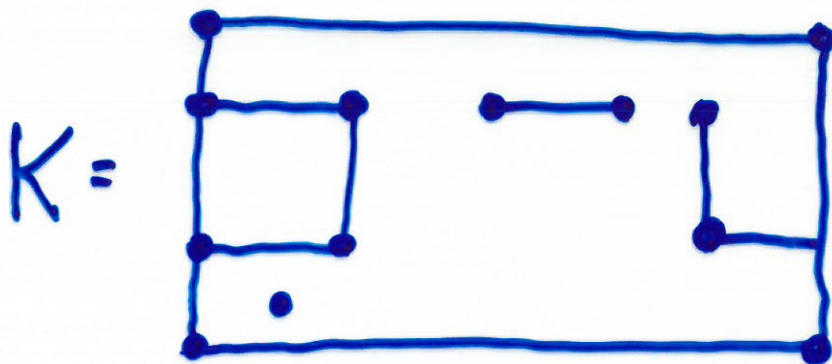
• IDENTITIES ARE

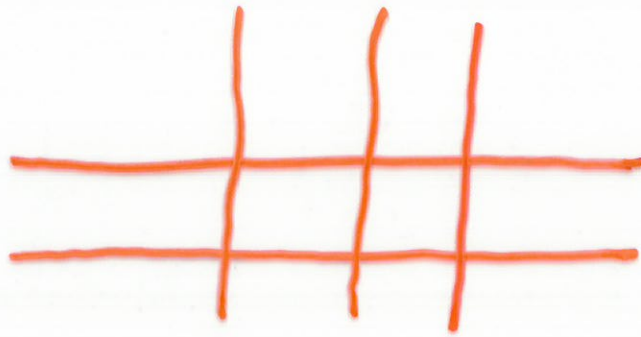


$$F: \begin{array}{ccc} [19] \xrightarrow{\cong} [2] \\ \downarrow \quad \downarrow \\ [2] \xrightarrow{\cong} [1] \end{array} \longrightarrow \mathcal{A}$$

WANT TO EXTEND F TO A DOUBLE

$$\text{FUNCTOR } F: \mathbb{RC} \longrightarrow \mathcal{A}$$







ADMISSIBLE SUBDIVISION : AT MOST ONE INTERIOR VERTEX INSIDE EACH SUBRECTANGLE.

$$[K] = ([K_{11}] \circ_h [K_{12}] \circ_h \dots \circ_h [K_{1n}]) \circ_v ([K_{21}] \circ_h \dots) \circ_v \dots$$

EACH SUBRECT IS ONE OF THE 19

BASIC ONES : $[K_{11}] =$  , $[K_{33}] =$ 

MUST HAVE

$$F[K] = (F[K_{11}] \circ_h F[K_{12}] \circ_h \dots) \circ_v (F[K_{21}] \circ_h \dots) \dots$$

- UNIQUE
- INDEPENDENT OF SUBDIVISION
- HOMOTOPY INVARIANT
- DOUBLE FUNCTOR

COLOURED COMPLEXES

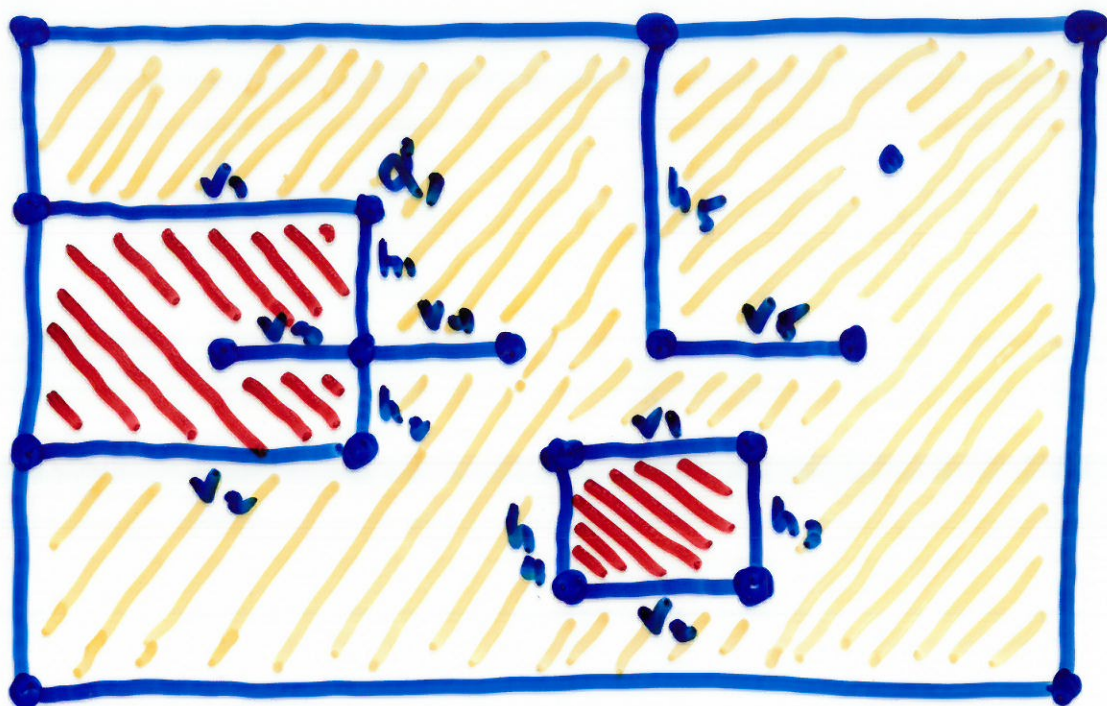
\mathbb{D} A DOUBLE GRAPH (REFLEXIVE)

$K = (R, P, H, V)$ A RECT COMPLEX

A **D-COLOURING** OF K IS A FN

$$\chi : R \longrightarrow \mathbb{D}_{11}$$

- EACH INTERIOR POINT OF R HAS A NBHD LIKE ONE OF THE 19 BASIC CELLS. THAT POINT MUST BE COLOURED BY A CELL OF \mathbb{D} OF SAME BOUNDARY TYPE.
- FURTHER, ALL POINTS OF THE NBHD ARE COLOURED WITH THE APPROPRIATE BOUNDARIES OF THAT CELL
- BOUNDARY POINTS OF R ARE COLOURED LIKE THEIR NBRS.

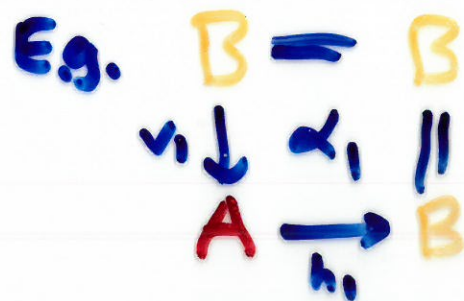


TWO OBJ: A B

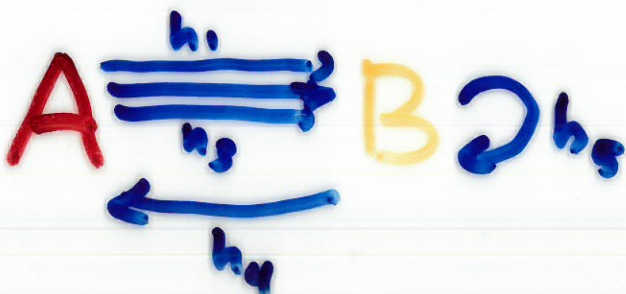
5 VERT ARR:





12 CELLS



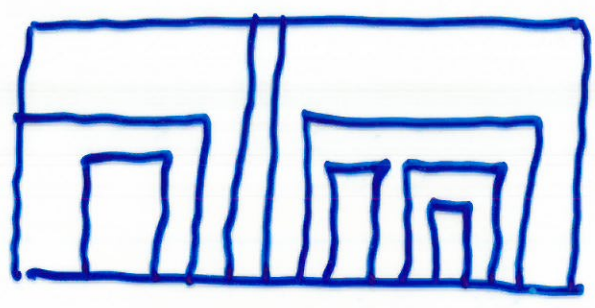
5 HOR ARR:



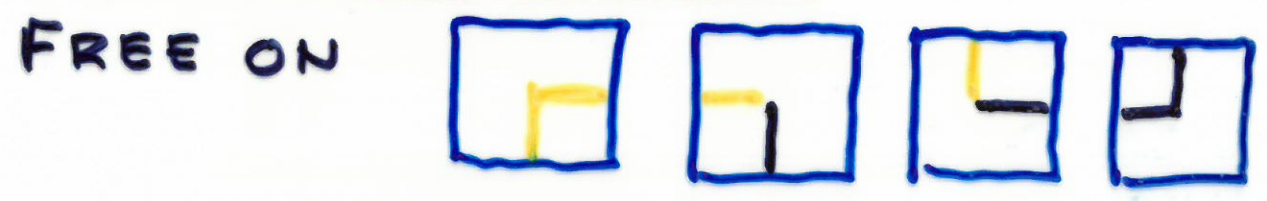
EXAMPLES

- THE CATALAN DOUBLE CATEGORY
FREE ON  

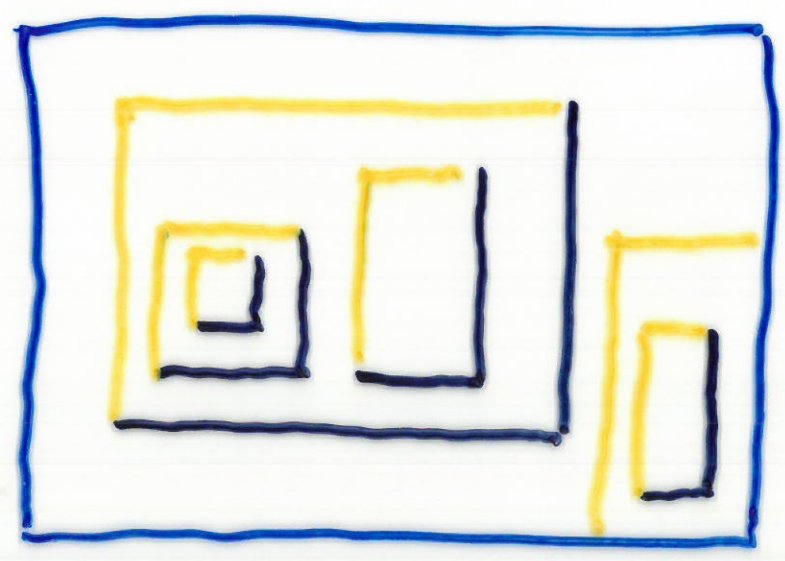
GENERAL CELL



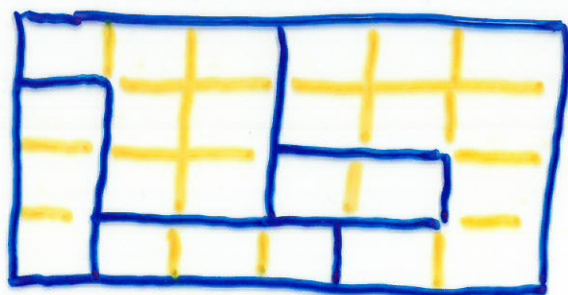
- STACKED RECTANGLES



GENERAL CELL



POLYOMINO TILINGS OF RECTANGLES



• 1 OBJ

• 2 HOR → →

• 2 VERT → →

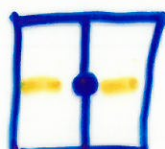
• 12 DOUB



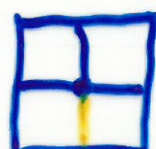
(1)



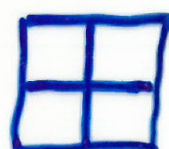
(4)



(2)



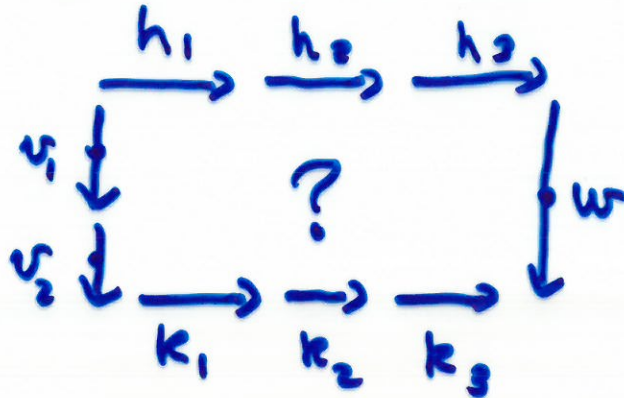
(4)



(1)

HOW COMPLEX ARE FREE DOUB CATS?

- ONE MEASURE : DOES THERE EXIST A CELL WITH A GIVEN BOUNDARY?



- FOR FREE **CATEGORY** ON A REFLEXIVE GRAPH WITH n VERTICES
 $\exists A \rightarrow B$ in FREE $\Leftrightarrow \exists$ PATH OF LENGTH n IN GRAPH. IF A IS INCIDENCE MATRIX OF GRAPH, ANS CAN BE READ OFF A^n .
- FOR FREE **DOUBLE CATEGORY**
 DIFFERENT STORY : THE WORD PROBLEM FOR GROUPS CAN BE MODELED IN ONE \Rightarrow **UNDECIDABLE**

WORD PROBLEM FOR GROUPS IS SPECIAL CASE OF WORD PROBLEM FOR CATS.

- G FINITE GRAPH

- R FINITE SET OF RELATIONS

$$r: \begin{array}{c} A_0 \xrightarrow{f_1} A_1 \xrightarrow{f_2} \dots \xrightarrow{f_m} A_m \\ \parallel \quad \parallel \quad \quad \quad \parallel \\ B_0 \xrightarrow{g_1} B_1 \xrightarrow{g_2} \dots \xrightarrow{g_m} B_m \\ \parallel \quad \parallel \quad \quad \quad \parallel \\ A \quad \quad \quad \quad \quad \quad \quad B \end{array}$$

$\langle G | R \rangle = \text{FREE CAT ON } G / \text{CONG GEN BY } R$

WHEN ARE 2 WORDS EQUIVALENT?

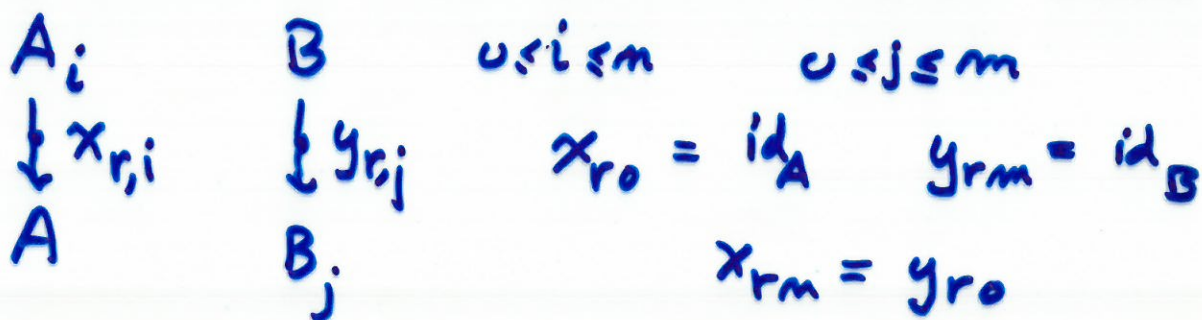
- CAN ASSUME R SYMMETRIC.

- CONSTRUCT A DOUBLE CATEGORY WHICH TURNS RELATIONS INTO CELLS.

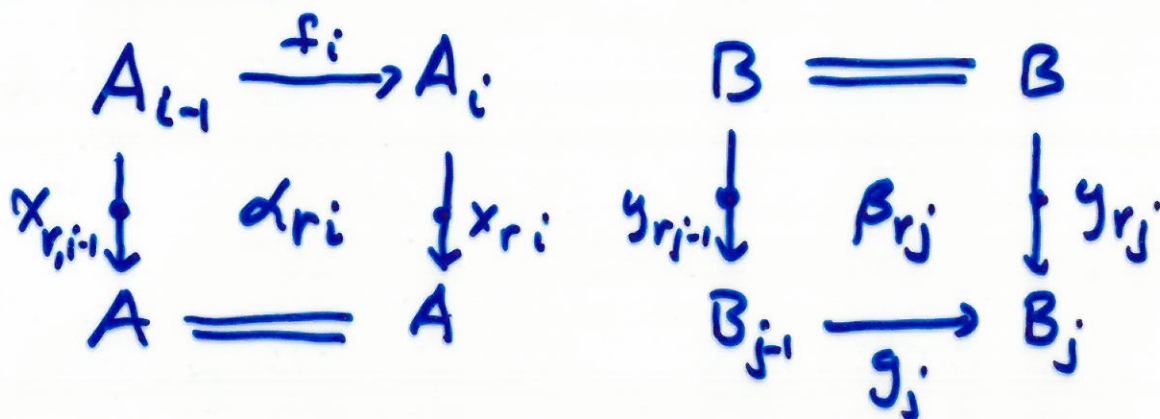
FREE

THE VERTICIES AND HORIZONTAL ARROWS OF ID ARE THOSE OF G .

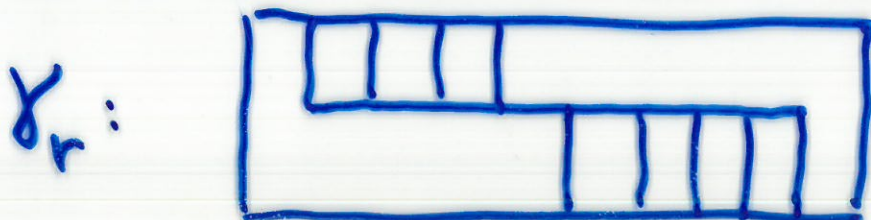
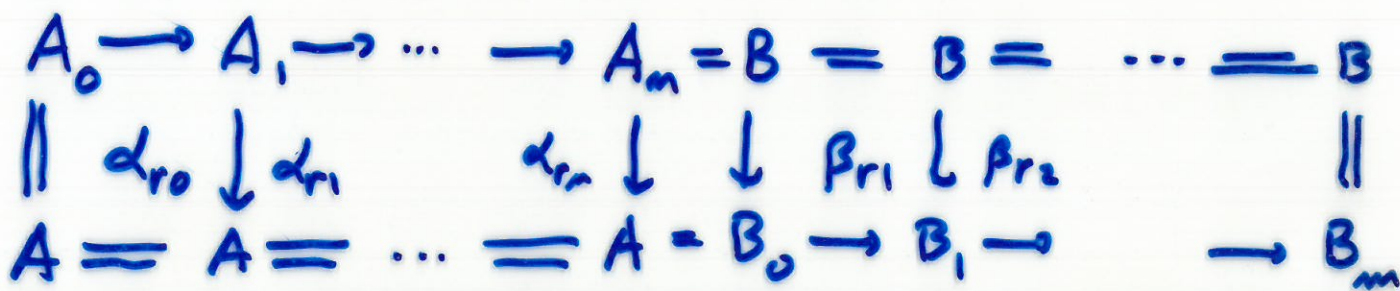
• FOR EACH $r \in \mathbb{R}$ DISTINCT VERT ARROWS



AND CELLS



• ONLY HOR COMPOSITS OF α 's & β 's IS WITH NEXT IN LINE



THM: GIVEN TWO PARALLEL WORDS
 $c_1 c_2 \dots c_k$ & $d_1 d_2 \dots d_l$, THEN THERE
 EXISTS A CELL δ WITH BOUNDARY

$$\begin{array}{ccccccc} C & \xrightarrow{\alpha_1} & C_1 & \xrightarrow{\alpha_2} & C_2 & \rightarrow \dots & \rightarrow C_{k-1} & \xrightarrow{\alpha_k} & D \\ \parallel & & & & & & & & \parallel \\ C & \xrightarrow{d_1} & D_1 & \xrightarrow{d_2} & D_2 & \rightarrow \dots & \rightarrow D_{l-1} & \xrightarrow{d_l} & D \end{array}$$

IN FREE(ID) IFF

$$c_1 c_2 \dots c_k = d_1 d_2 \dots d_l \text{ in } \langle G | R \rangle.$$

PROOF: (\Leftarrow) EASY

(\Rightarrow) GEOMETRIC + INDUCTION ON # α 's
 & β 's.

