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DOUBLE LIMITS

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LIMITS IN 2-CATEGORIES

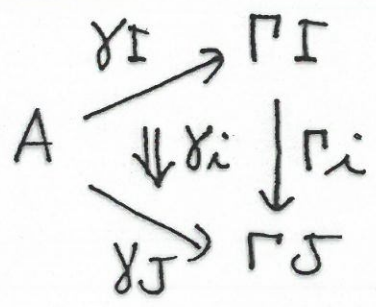
2- \lim , PSEUDO- \lim , LAX \lim , OP-LAX \lim

BI- \lim , WEIGHTED \lim , FLEXIBLE \lim

$\Gamma : I \rightarrow A$ 2-FUNCTOR

$\lim \Gamma$ IS A UNIVERSAL CONE ON Γ

CONE



UNIVERSAL

$A(X, A) \rightarrow \underline{CONE}(X, \Gamma)$

- γ_i COMPOSE
- COULD BE IDENTITIES, ISOS, MORPHISMS, OP-MOR.

- ISO OF CATS
- EQUIV OF CATS

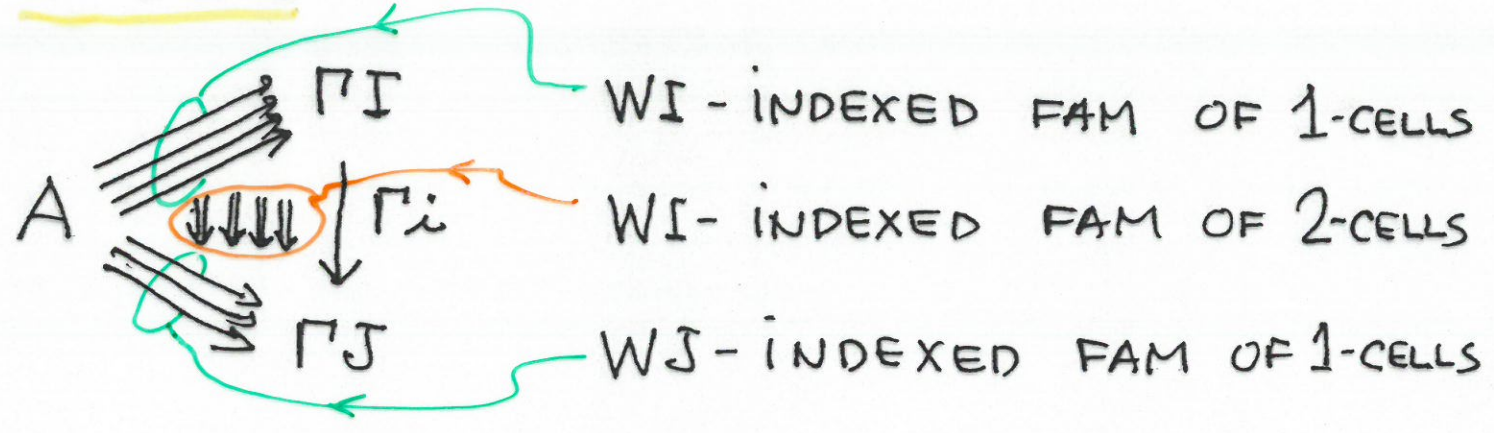
NOT QUITE GENERAL ENOUGH

INTRODUCE WEIGHTED $2\text{-}\varprojlim$

COMES FROM \underline{V} -CATEGORIES

WEIGHT: $W: \mathbb{I} \rightarrow \text{CAT}$ 2-FUNCT.

W-CONE



i.e. $W \xrightarrow{\gamma} \underline{A}(A, \Gamma -)$ (2-NAT, LAX-NAT, ...)

$\varprojlim_W \Gamma$ IS UNIVERSAL W-CONE.

PSEUDO-, LAX-, OP-LAX- \varprojlim ALL

CAN BE REDUCED TO WEIGHTED $2\text{-}\varprojlim$

$2\text{-}\varprojlim$ ARE TOO GENERAL

INTRODUCE FLEXIBLE \varprojlim

IN ORDER TO HANDLE ALL THESE
CONCEPTS, NEED SOME THEORY
TO GUIDE US.

THE THEORY USED IS THAT OF
V-CATEGORIES FOR V = CAT

PROPOSAL: USE CAT-INDEXED CAT THY

PROGRAM: ORGANIZE & EXTEND
THE THEORY OF 2-CATS USING
INDEXED CATS

REFINE THE THEORY
OF INDEXED CATS BY STUDYING
2-CATEGORIES

CATEGORY OF ELEMENTS

5

$$F: \underline{A}^{\text{op}} \rightarrow \underline{\text{SET}} \rightsquigarrow \begin{array}{c} \text{EL}(F) \\ \downarrow P \\ \underline{A} \end{array}$$

- F REPRESENTABLE $\Leftrightarrow \text{EL}(F)$ HAS TERM. OBJ.
- $F = \underline{A}(-, A) \Rightarrow \text{EL}(F) = \underline{A}/A$
- F FLAT $\Leftrightarrow \text{EL}(F)$ IS FILTERED
- $\varinjlim_F \Phi \cong \varinjlim \Phi P$

FOR 2-CATS - MANY CHOICES

$F: \underline{A} \rightarrow \underline{\text{CAT}}$ 2-FUNCT

$\text{EL}(F)$ OBJ. $(A, X \in \text{FA})$

MORPH. $(A, X) \rightarrow (B, Y)$

$$a: A \rightarrow B$$

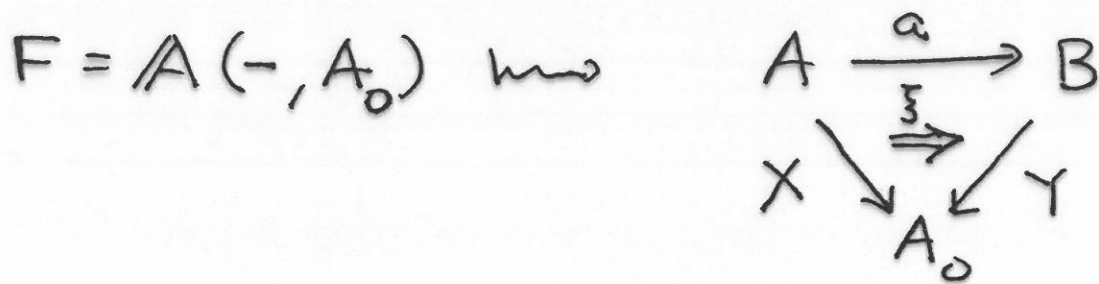
$$F(a)(Y) = X$$

$$F(a)(Y) \xrightarrow{\cong} X$$

$$X \rightarrow F(a)(Y)$$

$$F(a)(Y) \rightarrow X$$

2-CELLS ARE $A \Downarrow B$ + COMPATIBILITY



$\{$ EQUALITY, ISO, MORPH, OP-MORPH.

EACH GIVES A USEFUL CONCEPT

BUT NONE IS ADEQUATE FOR

EVERYTHING, E.G. REPRESENTABILITY

THEOREM. OR TRANSFORMING

WEIGHTED \varinjlim INTO UNWEIGHTED

ONES.

THE THEORY OF INDEXED CATEGORIES

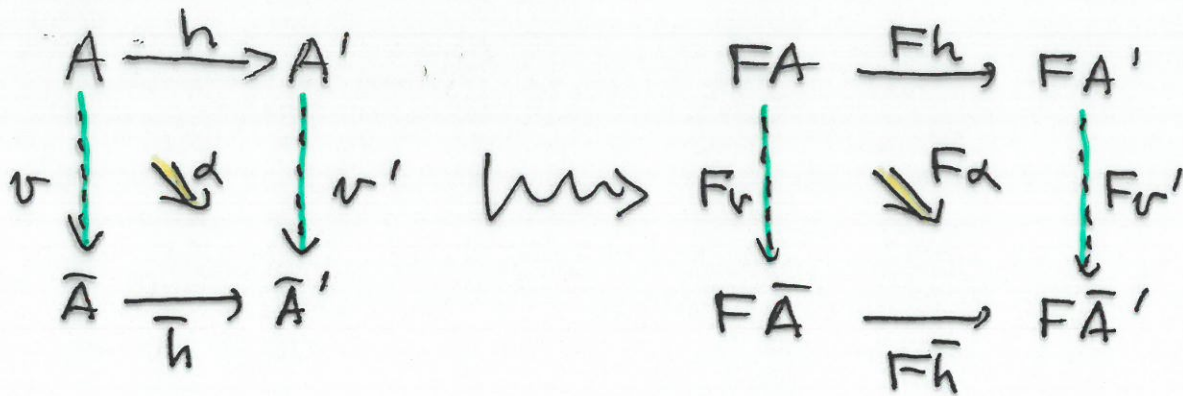
TELLS US THAT LIMITS SHOULD BE

TAKEN OVER CATEGORY OBJECTS IN

THE BASE, I.E. DOUBLE CATEGORIES.

Ex3: DOUB

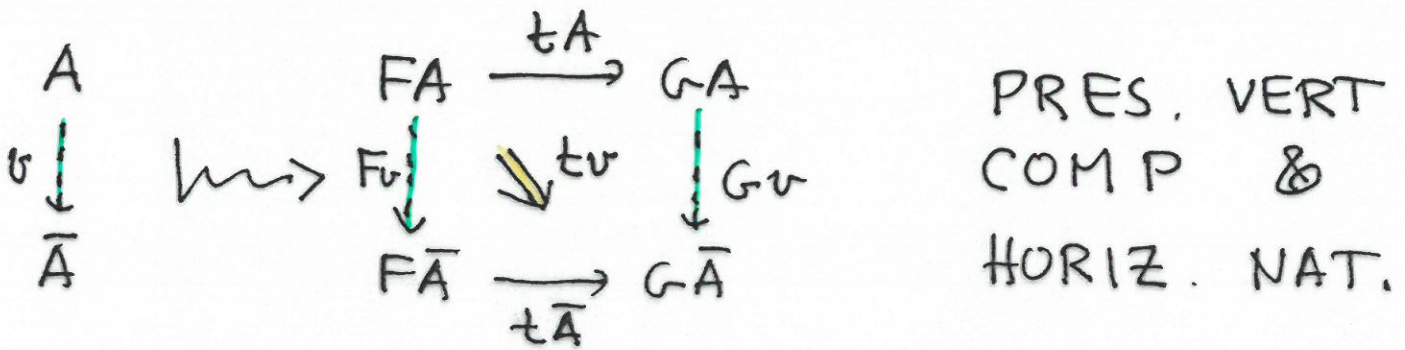
DOUBLE FUNCTORS $F: \mathbb{D} \rightarrow \mathbb{E}$



PRESERVES ALL COMPOSITIONS.

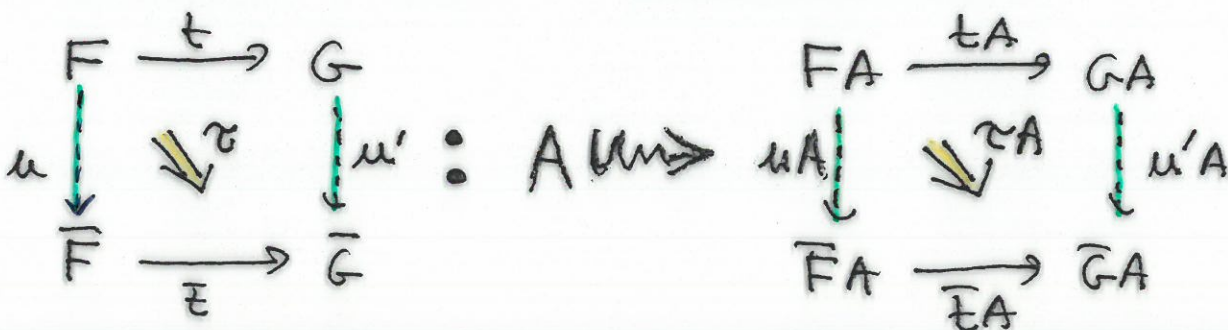
HORIZONTAL NATURAL TRANSFORMATIONS

$t: F \rightarrow G$



VERT. NAT. TRANSF. SIMILAR

DOUBLE NAT. TRANSFS.



CONSTRUCTIONS ON DOUBLE CATEGORIES

DOUB. CAT \rightsquigarrow CAT INDEXED CAT

$\mathbb{D} \rightsquigarrow \underline{\underline{D}} : \underline{\underline{D}}^{\underline{\underline{C}}}$ OBJ. $\underline{\underline{C}} \rightarrow \underline{\underline{D}}_0$
 MOR. $\underline{\underline{C}} \rightarrow \underline{\underline{D}}_1$

CAN RECOVER \mathbb{D}

$\underline{\underline{D}}^1 \rightsquigarrow$ HORIZ. (LEADING ASPECT)

$\underline{\underline{D}}^2 \rightsquigarrow$ VERT & DOUB

$\underline{\underline{D}}^3 \rightsquigarrow$ VERT COMP

USE THIS PROCESS TO TRANSLATE CONCEPTS WELL UNDERSTOOD FOR CAT IND. CATS TO DOUBLE CATS.

1. FUNCTOR CATEGORIES

$\mathbb{D}^{\mathbb{E}}$ DOUBLE FUNCTORS, HORIZ.
 VERT. & DOUBLE NAT. TRANSFS

FOR \mathbb{D} & \mathbb{E} 2-CATEGORIES $\mathbb{D}^{\mathbb{E}}$

IS NOT GENERALLY A 2-CAT

E.G. $\mathbb{D} = \text{CAT}$, $\mathbb{E} = \mathbb{2}$ (VERT. DISC.)

$\text{CAT}^{\mathbb{2}}$

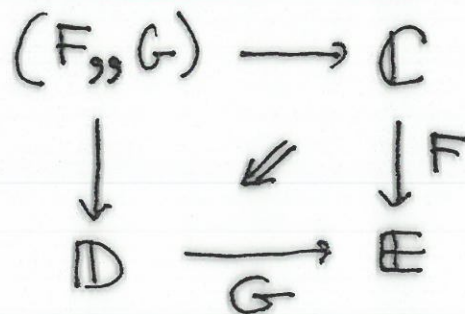
OBJ. $\underline{A} \xrightarrow{F} \underline{B}$

HORIZ. $\begin{array}{ccc} \underline{A} & \xrightarrow{H} & \underline{A'} \\ F \downarrow & \equiv & \downarrow F' \\ \underline{B} & \xrightarrow{K} & \underline{B'} \end{array}$

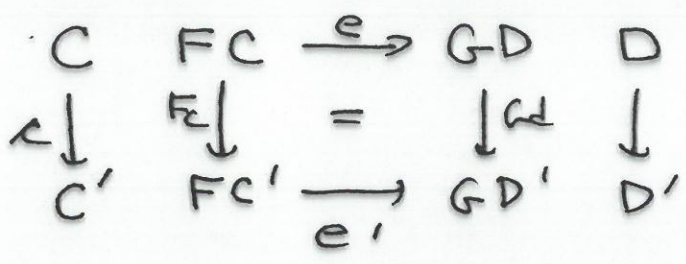
VERT. $\underline{A} \begin{array}{c} \xrightarrow{F} \\ \Downarrow t \\ \xrightarrow{G} \end{array} \underline{B}$

DOUB. $\begin{array}{ccc} \underline{A} & \xrightarrow{H} & \underline{A'} \\ F \left(\begin{array}{c} \rightarrow \\ \downarrow \\ \rightarrow \end{array} \right) G & \equiv & F' \left(\begin{array}{c} \rightarrow \\ \downarrow \\ \rightarrow \end{array} \right) G' \\ \underline{B} & \xrightarrow{K} & \underline{B'} \end{array}$

2. COMMA CATEGORIES (HORIZ.)

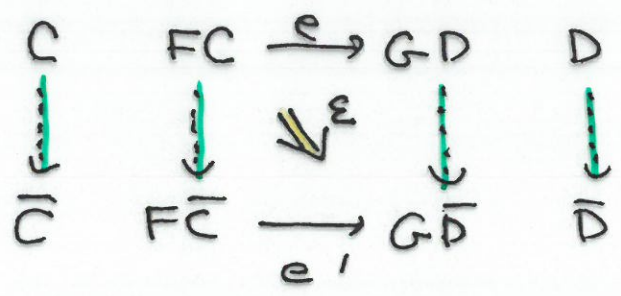


HORIZ.

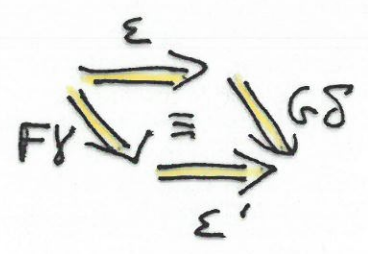
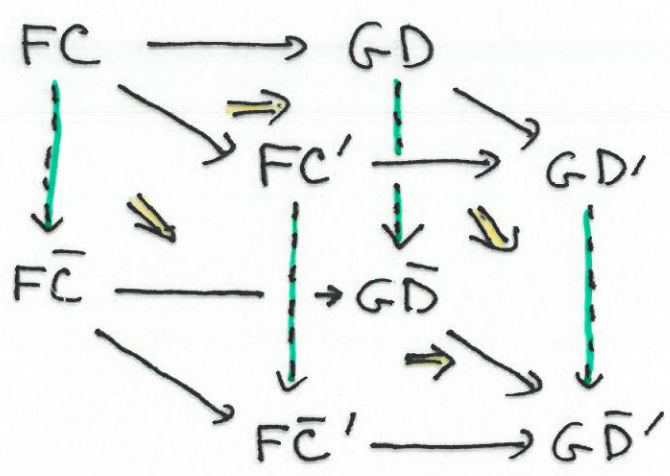


(ALL HORIZ.)

VERT.



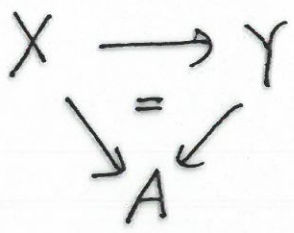
DOUB.



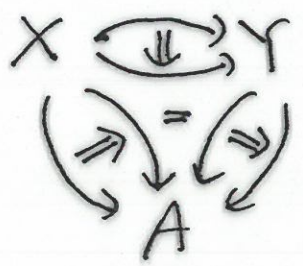
3. DOUBLE SLICE

A // A

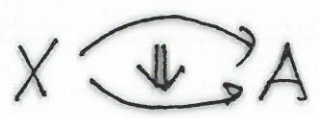
HORIZ.



DOUB



VERT.



4. "CATEGORY OF ELEMENTS"

$$F: \mathcal{A}^{op} \rightarrow \text{CAT} \quad \text{2-FUNCTOR}$$

$$\text{EL}(F) \quad \text{OBJ.} \quad (A, X \in FA)$$

$$\text{HORIZ.} \quad \frac{(A, X \in FA) \longrightarrow (B, Y \in FB)}{A \xrightarrow{a} B \quad F(a)(Y) = X}$$

$$\text{VERT.} \quad \frac{(A, X \in FA) \longrightarrow (A, X' \in FA)}{x: X \longrightarrow X'}$$

$$\text{DOUB.} \quad \begin{array}{ccc} (A, X) & \xrightarrow{(a, Y)} & (B, Y) \\ \downarrow & \searrow (a, y) & \downarrow \\ (A, x) & & (B, y) \\ (A, X') & \xrightarrow{(b, Y')} & (B, Y') \end{array}$$

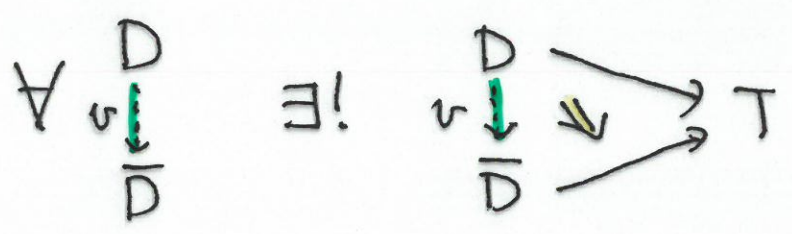
WHERE $A \begin{array}{c} \xrightarrow{a} \\ \Downarrow \alpha \\ \xrightarrow{b} \end{array} B$ AND

$$x = F(\alpha)(Y') \cdot F(a)(y) = F(b)(y) \cdot F(\alpha)(Y)$$



NAT OF F(a)

A TERMINAL OBJECT IN A DOUBLE
 CAT \mathbb{D} IS T SUCH THAT
 $\forall D \exists ! D \rightarrow T$ AND



THEOREM : $F : \mathcal{A}^{op} \rightarrow \text{CAT}$ IS
 REPRESENTABLE $\iff \text{EL}(F)$ HAS TERM. OBJ.

IF $F = \mathcal{A}(-, A)$ THEN $\text{EL}(F) = \mathcal{A} \parallel A$
 AND $1_A : A \rightarrow A$ IS TERMINAL.

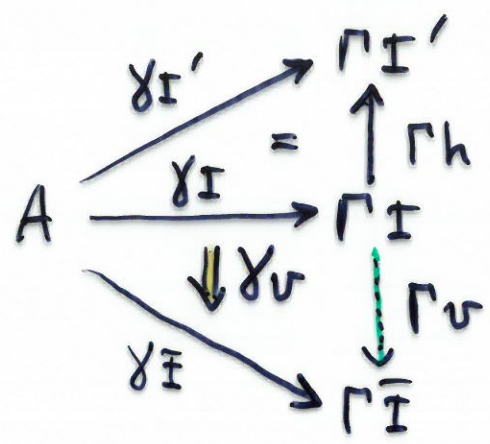
DOUBLE CONES

$\Gamma : \mathbb{I} \rightarrow \mathbb{D}$ A DOUBLE DIAGRAM.

A DOUBLE CONE IS A HORIZ.

NAT TRANSF $\text{CONST}(A) \rightarrow \Gamma$
 FOR A AN OBJ. OF \mathbb{D}

I.E.

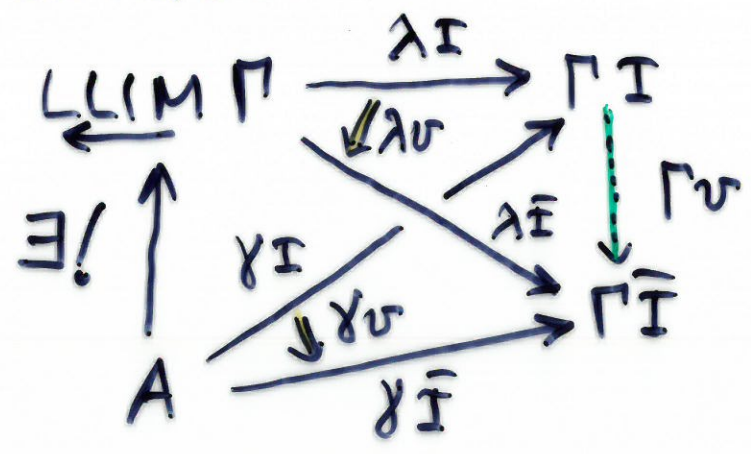


γ_I HORIZ. NAT.
 γ_v COMPOSE VERT.
 HORIZ. NAT.
 W.R.T. DOUBLE
 MAPS

THE DOUBLE LIMIT $\varprojlim \Pi$
 IS A UNIVERSAL SUCH CONE

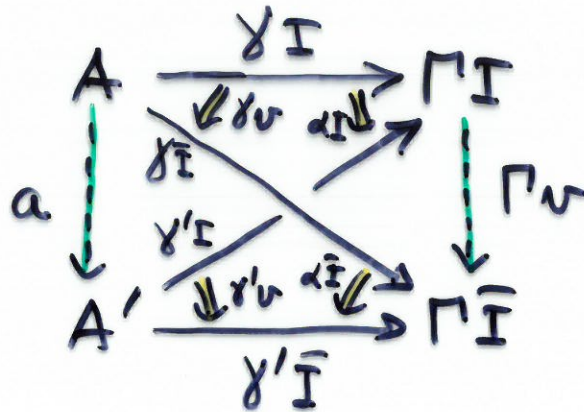
UNIVERSALITY

ANY CONE SUCH AS γ ABOVE
 FACTORS UNIQUELY

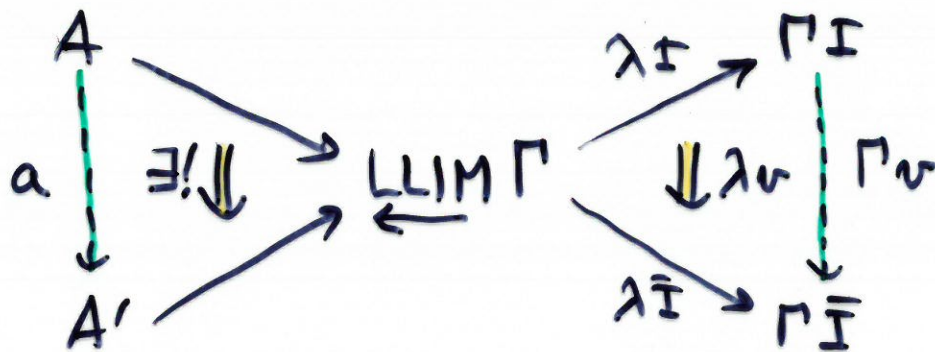


AND...

ANY "DOUBLE CONE"



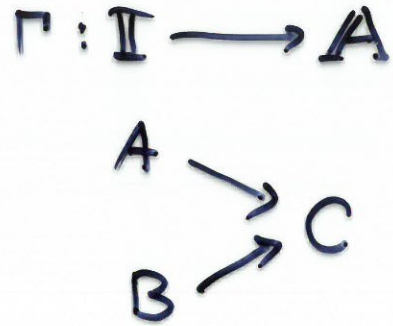
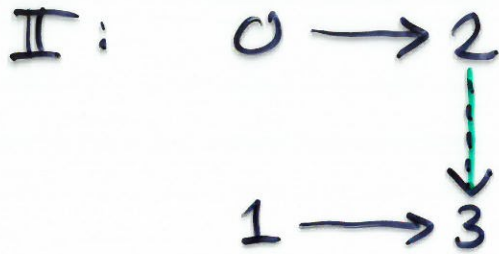
FACTORS UNIQUELY AS



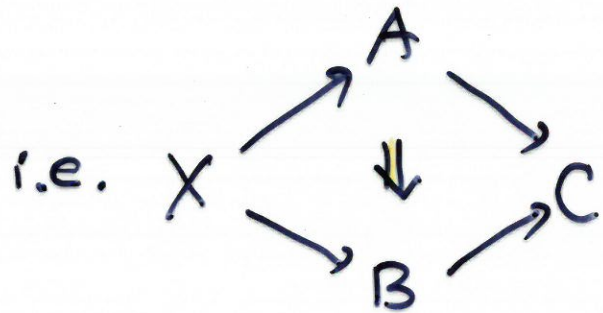
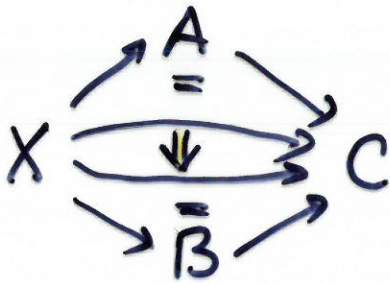
THIS CONCEPT CAN BE USED IN A 2-CAT. HERE THE Γv WOULD ALL BE IDENTITIES BUT NOT λv .

Ex 0. IF $\Gamma: \emptyset \rightarrow \mathcal{A}$ IS THE EMPTY DIAGRAM, THEN $\underline{\text{LLIM}} \Gamma$ IS A TERMINAL OBJECT, AS ABOVE.

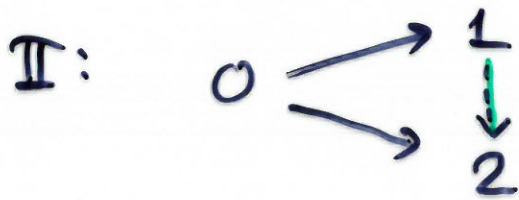
Ex 1. COMMA OBJECTS



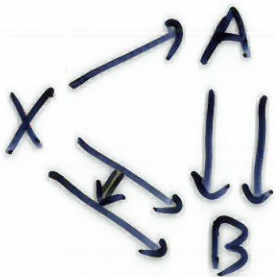
CONE IS



Ex 2. INSERTERS



CONE IS



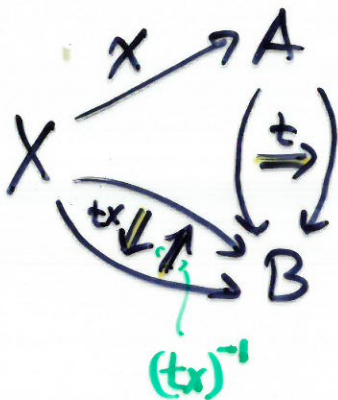
Ex 3. INVERTERS



$$\Gamma : \mathbb{I} \rightarrow A$$

$$A \xrightarrow{\downarrow} B$$

COME IS



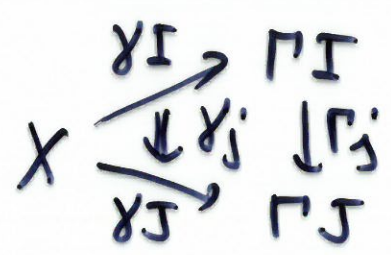
$$\text{I.E. } X \xrightarrow{x} A \xrightarrow{\downarrow t} B$$

S.T. tx INV.

EX 4. LAX LIMITS

$\Gamma: \underline{I} \rightarrow \mathcal{A}$ 2-FUNCT (\underline{I} A CAT)

LAX CONE



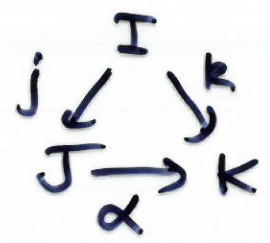
THE γ_j
COMPOSE

CONSTRUCT A DOUBLE CAT

$\mathcal{B}\underline{I}$ OBJ \underline{I} HORIZ

$\downarrow j$

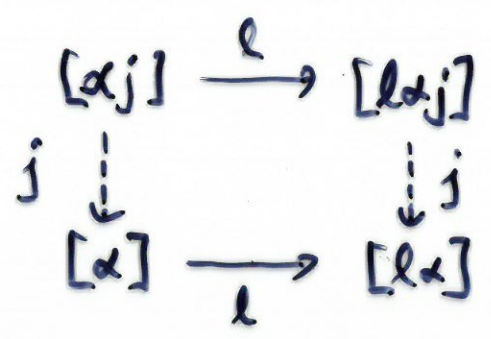
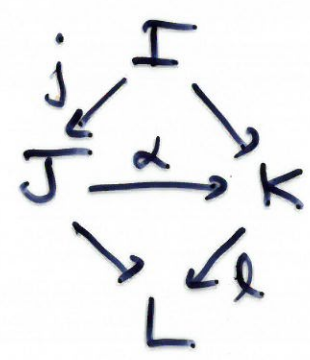
\underline{J}



VERT $\underline{I} \xrightarrow{\alpha} \underline{I}'$ DOUB

$j \swarrow \quad \searrow j'$

\underline{J}



$D_1: \mathcal{B}\underline{I} \rightarrow \underline{I}$ TAKE CODOMAIN

THEN $\text{LAX } \Gamma \cong \text{LLIM } \Gamma D_1$

THEOREM : LET $W: \mathbb{I} \rightarrow \text{CAT}$ BE
 A 2-FUNCTOR AND $P: \mathbb{E}L(W) \rightarrow \mathbb{I}$
 THE CANONICAL PROJECTION. THEN
 FOR EVERY $\Gamma: \mathbb{I} \rightarrow \mathbb{A}$ (2-FUNCTOR)

$$\lim_{\leftarrow} W \Gamma \cong \lim_{\leftarrow} \Gamma P$$

EX: PSEUDO & LAX LIM ARE LLIM'S

PROPOSITION: EVERY DOUBLE LIMIT
 IS AN ITERATED WEIGHTED 2-LIM.
 IN PARTICULAR 2-COMPLETENESS
 \Rightarrow DOUBLE COMPLETENESS.

"PROOF" FOR INDEXED CATS
 LIM IS EQUALIZER OF π .

EQUALIZERS ARE 2-LIM'S.

π IS \lim_{\leftarrow} WHEN \mathbb{I} IS
 HORIZONTALLY DISCRETE.

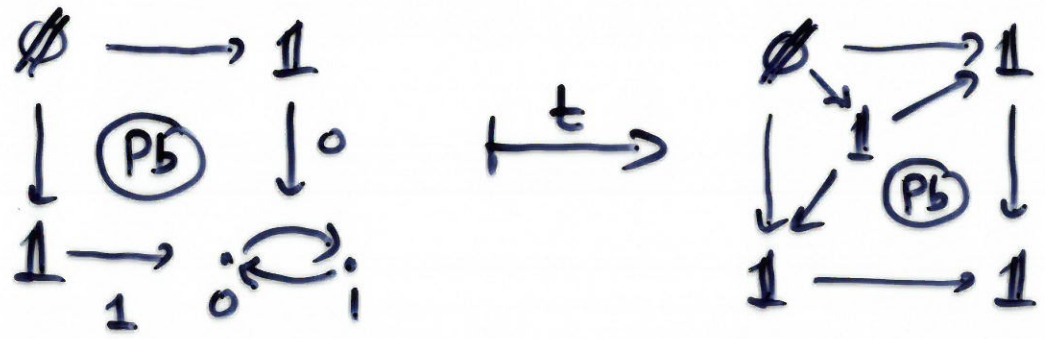
$\Gamma: \mathbb{I} \rightarrow \mathbb{A}$ (2-CAT) IS CONSTANT
 ON COMPONENTS.

$$\lim_{\leftarrow} \Gamma = \prod_{[I] \in \pi_0 \mathbb{I}} \Gamma(I)^{[I]}$$

PERSISTENT DOUBLE LIMITS

DEF: \mathbb{I} LLIM ARE PERSISTENT
 IF FOR ANY TWO DIAGRAMS
 $\Gamma, \Phi: \mathbb{I} \rightarrow \mathcal{A}$ AND ANY HORIZ.
 NAT. TRANSF. $t: \Gamma \rightarrow \Phi$
 SUCH THAT $\forall I, t_I: \Gamma_I \rightarrow \Phi_I$
 IS AN EQUIVALENCE, IT FOLLOWS
 THAT LLIM $\Gamma \rightarrow$ LLIM Φ IS
 AN EQUIVALENCE.

E.G. PULLBACKS ARE NOT PERSISTENT

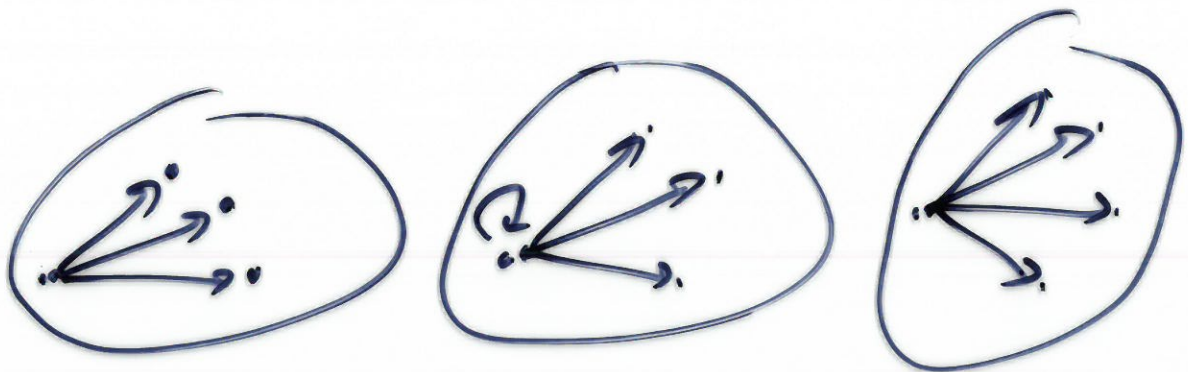


COMMA OBJECTS ARE.

THEOREM: \mathbb{I} $\xleftarrow{\text{LLIM}}$ ARE
 PERSISTENT IFF THE HORIZONTAL
 PART OF \mathbb{I} IS COMPONENTWISE
 CONTRACTIBLE.

THIS MEANS THAT $\exists R: \pi_0 \mathbb{I}_h \longrightarrow \mathbb{I}_h$
 AND $\exists t: RQ \longrightarrow 1_{\mathbb{I}_h}$ SUCH THAT
 $QR = 1_{\pi_0 \mathbb{I}_h}$, WHERE $Q: \mathbb{I}_h \longrightarrow \pi_0 \mathbb{I}_h$
 IS THE QUOTIENT MAP.

MORE CONCRETELY: EACH COMPONENT
 OF \mathbb{I}_h HAS A NATURAL WEAK
 INITIAL OBJECT



NOTE: FLEXIBLE \implies PERSISTENT