

MATH/CSCI 2112, DISCRETE STRUCTURES I, FALL 2005
Handout 6: Rules for Program Correctness Proofs

Precondition strengthening

$$\frac{P_1 \Rightarrow P \quad \{P\} C \{Q\}}{\{P_1\} C \{Q\}}$$

Postcondition weakening

$$\frac{\{P\} C \{Q\} \quad Q \Rightarrow Q_1}{\{P\} C \{Q_1\}}$$

Conjunction rule

$$\frac{\frac{\{P_1\} C \{Q_1\} \quad \{P_2\} C \{Q_2\}}{\{P_1 \wedge P_2\} C \{Q_1 \wedge Q_2\}}}{\{P_1 \wedge P_2\} C \{Q_1 \wedge Q_2\}}$$

Disjunction rule

$$\frac{\frac{\{P_1\} C \{Q_1\} \quad \{P_2\} C \{Q_2\}}{\{P_1 \vee P_2\} C \{Q_1 \vee Q_2\}}}{\{P_1 \vee P_2\} C \{Q_1 \vee Q_2\}}$$

Sequencing rule (concatenation)

$$\frac{\frac{\{P\} C_1 \{Q\} \quad \{Q\} C_2 \{R\}}{\{P\} C_1; C_2 \{R\}}}{\{P\} C_1; C_2 \{R\}}$$

Modified concatenation rule

$$\frac{\frac{\{P\} C_1 \{Q\} \quad Q \Rightarrow R \quad \{R\} C_2 \{S\}}{\{P\} C_1; C_2 \{S\}}}{\{P\} C_1; C_2 \{S\}}$$

Assignment rule for variables

$$\frac{}{\{Q[E/x]\} x := E \{Q\}}$$

Modified assignment rule for variables

$$\frac{}{\{E \text{ well-defined, } Q[E/x]\} x := E \{Q\}}$$

If-then-else rule, forward

$$\frac{\frac{\{P \wedge B\} C_1 \{Q\} \quad \{P \wedge \sim B\} C_2 \{R\}}{\{P\} \text{ if } B \text{ then } C_1 \text{ else } C_2 \{Q \vee R\}}}{\{P\} \text{ if } B \text{ then } C_1 \text{ else } C_2 \{Q \vee R\}}$$

If-then-else rule, backward

$$\frac{\frac{\{P\} C_1 \{Q\} \quad \{R\} C_2 \{Q\}}{\{(P \wedge B) \vee (R \wedge \sim B)\} \text{ if } B \text{ then } C_1 \text{ else } C_2 \{Q\}}}{\{(P \wedge B) \vee (R \wedge \sim B)\} \text{ if } B \text{ then } C_1 \text{ else } C_2 \{Q\}}$$

If-then-else rule, both ways

$$\frac{\frac{\{P \wedge B\} C_1 \{Q\} \quad \{P \wedge \sim B\} C_2 \{Q\}}{\{P\} \text{ if } B \text{ then } C_1 \text{ else } C_2 \{Q\}}}{\{P\} \text{ if } B \text{ then } C_1 \text{ else } C_2 \{Q\}}$$

If-then rule, both ways

$$\frac{\frac{\{P \wedge B\} C_1 \{Q\} \quad P \wedge \sim B \Rightarrow Q}{\{P\} \text{ if } B \text{ then } C_1 \{Q\}}}{\{P\} \text{ if } B \text{ then } C_1 \{Q\}}$$

While rule

$$\frac{\frac{\forall k. (\{I(k) \wedge B\} C \{I(k+1)\}) \quad P \Rightarrow I(0) \quad I(k) \wedge \sim B \Rightarrow Q \quad \exists k. (I(k) \Rightarrow \sim B)}{\{P\} \text{ while } B \text{ do } C \{Q\}}}{\{P\} \text{ while } B \text{ do } C \{Q\}}$$

Problem 1. In each of parts (a)–(e), find the weakest precondition P that makes the given Hoare triple correct. Simplify your answers (logically and algebraically) where applicable.

- (a) $\{P\} x := 5 - x \{n = x + 7\}$
- (b) $\{P\} a := -a; n := n + a \{n \neq a \wedge n \neq 0\}$
- (c) $\{P\} a := a + b; b := a - b; a := a - b \{a \geq b + 2\}$
- (d) $\{P\} \text{ if } x > y \text{ then } x := x + 2 \text{ else } y := y - 2 \{x \neq y\}$
- (e) $\{P\} \text{ if } x > y \text{ then } z := x \text{ else } z := y \{z > y\}$