

FIBONACCI NUMBERS ARE NOT CONTEXT-FREE

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The Fibonacci numbers, given by the recurrence relation

$$F(n+2) = F(n+1) + F(n), \quad F(1) = 1, \quad F(2) = 1,$$

are considered to be written in base b , so "trailing zeros" correspond exactly to "factors of b ." From [4], Theorem 5, page 527, it follows that, for any prime p , there exists n s.t. $F(k \times n) \equiv 0 \pmod{p^i}$ for positive i and k . The existence of j s.t. $F(j) \equiv 0 \pmod{b^i}$, for arbitrary positive b , follows by applying the above to the prime factoring of b and choosing j to be the least common multiple of the n . Thus, in any base, there exist Fibonacci numbers with arbitrarily many trailing zeros.

In the proof of this same theorem [4], it is established for any prime p that, if $F(n)$ is the first term $\equiv 0 \pmod{p^e}$ but $\not\equiv 0 \pmod{p^{e+1}}$, then $F(p \times n)$ is the first term $\equiv 0 \pmod{p^{e+1}}$, also $F(p \times n) \not\equiv 0 \pmod{p^{e+2}}$.

This establishes, for each prime base p , a lower bound on n which increases exponentially with the number of trailing zeros in $F(n)$ base p . This bound generalizes to composite bases because when $F(n)$ has e trailing zeros in base b it must also have e trailing zeros in all bases p , where p is a prime factor of b . Specifically, there is some constant k such that, for all sufficiently large n ,

$$TZ(F(n)) < k \times \log(n),$$

where $TZ(x)$ is the number of trailing zeros in x .

Since the Fibonacci sequence is asymptotically exponential, there is some constant c s.t. $n < c \times |F(n)|$, where $|F(n)|$ denotes the *length* of $F(n)$ as a string, i.e., the number of digits in $F(n)$ in base b . Combining these, and adjusting k to also account for c , gives

$$(1) \quad TZ(F(n)) < k \times \log(|F(n)|).$$

These facts can be used to show that the Fibonacci numbers do not form a *context-free* set. A set of strings is context-free iff it is the set generated by some context-free grammar or, equivalently, a set of strings is context-free iff it is the set recognized by some pushdown automaton. Ogden's Lemma, stated below, gives a property true of all context-free sets, and is used in Lemma 1 to show a set of strings closely related to the Fibonacci numbers to be not context-free.

Ogden's Lemma [2]: Let Q be a context-free set. Then there is a constant j such that, if α is any string in Q and we mark any j or more positions of α "distinguished," then we can write $\alpha = uvwxy$, such that:

- 1) v and x together have at least one distinguished position,
- 2) vwx has at most n distinguished positions, and
- 3) for all $i \geq 0$, uv^iwx^iy is in Q .

Lemma 1: Let Q be the set of strings such that the members of Q are the Fibonacci numbers written in base b with a new symbol "#" inserted immediately following the last nonzero digit. The set Q is not context-free.

Proof: The proof is by contradiction. Assume that Q is context-free.

Let j be the number of "distinguished" positions required for Ogden's Lemma (see [2] for a description of Ogden's Lemma). Since we know there are

Fibonacci numbers with arbitrarily many trailing zeros, let α be a member of Q corresponding to a Fibonacci number with at least j trailing zeros. The trailing zeros, which follow the "#," are used as the distinguished positions for purposes of Ogden's Lemma.

Applying Ogden's Lemma, α may be partitioned as follows:

$$\alpha = uvwxy,$$

where x contains at least one of the trailing zeros. Further, for all $i \geq 0$, $\beta_i = uv^iwx^i y$ must also be in the set Q , and thus correspond to some Fibonacci number satisfying (1).

If x contained the "#," then clearly β_2 would contain two "#" symbols and, thus, could not be a member of Q . Therefore, x contains only "0"s, so β_i has at least $j + i - 1$ trailing zeros.

Since v and x together can be no longer than α , then β_i can be no more than i times as long as α : So $|\beta_i| \leq i \times |\alpha|$. Applying (1) to these bounds gives:

$$j + i - 1 < k \times \log(i \times |\alpha|).$$

Choosing $i = 2k^2|\alpha| + 1$ produces a contradiction. \square

Theorem: For all integers $b \geq 2$, the set of Fibonacci numbers in base b , considered as strings over the alphabet $0, 1, \dots, b - 1$, is not context-free.

Proof: Assume M is a pushdown automaton (PDA) recognizing the set of Fibonacci numbers. We modify the finite control to give another PDA M' , recognizing the set Q , thus contradicting Lemma 4. An informal description of M' follows.

M' contains a copy of the machine M , plus additional logic in the finite control to filter the input and pass it to this internal copy of M . M' accepts only when this internal M accepts the string passed to it.

behaves as follows:

- M' rejects if the input does not contain exactly one "#," if the "#" does not immediately follow a nonzero digit, or if there are any nonzero digits following the "#." Otherwise, M' accepts if and only if its internal simulation of M accepts.
- When M' reads a digit (any symbol except "#") from the input, it passes that digit to M . The "#" symbol, having been checked as above, is otherwise ignored and is not passed to M .

By the above rules, if M' accepts, then the input must be a Fibonacci number with a "#" inserted following the last nonzero digit. Thus, the input is in the set Q .

Conversely, if the input is in the set Q , then M' will pass the Fibonacci number to M and thus accept.

Therefore, M' accepts the set Q , a contradiction by Lemma 4; hence, the set of Fibonacci numbers is not context-free. \square

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