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1. RATIONAL CHEBYSHEV APPROXIMATIONS OF ANALYTIC FUNCTIONS

We proceed to establish the main result of this paper: a general procedure to obtain rational Chevyshev approximations of analytic functions. Let f(z) be analytic at z_0 . Then, by composition, $g(z) = f(\cos z + z_0 - 1)$ is analytic at the origin. Hence, we can write

$$g(z) = f(\cos z + z_0 - 1) = \sum_{n=0}^{\infty} g^{(2n)}(0) \frac{z^{2n}}{(2n)!}.$$
 (1.1)

If an explicit expansion of $f(\cos z + z_0 - 1)$ is not available, then successive coefficients in (1.1) are found directly from the formula for Maclaurin expansions, i.e., by simply calculating successive derivatives of (1.1) and setting z = 0. To wit,

$$g(0) = f(z_0), (1.2)$$

$$g''(0) = -f'(z_0), \tag{1.3}$$

$$g^{(iv)}(0) = 3f''(z_0) + f'(z_0), \tag{1.4}$$

$$g^{(\mathrm{vi})}(0) = -15f^{\prime\prime\prime}(z_0) - 15f^{\prime\prime}(z_0) - f^{\prime}(z_0), \tag{1.5}$$

$$g^{(\text{viii})}(0) = 105f^{(\text{iv})}(z_0) + 210f^{\prime\prime\prime}(z_0) + 63f^{\prime\prime}(z_0) + f^{\prime}(z_0), \tag{1.6}$$

$$g^{(x)}(0) = -945f^{(v)}(z_0) - 3150f^{(iv)}(z_0) - 2205f^{\prime\prime\prime}(z_0) - 255f^{\prime\prime}(z_0) - f^{\prime}(z_0),$$
(1.7)

$$g^{(\text{xii})}(0) = 10395 f^{(\text{vi})}(z_0) + 51975 f^{(\text{v})}(z_0) + 65835 f^{(\text{iv})}(z_0) + 21120 f^{'''}(z_0) + 1023 f^{''}(z_0) + f^{\prime}(z_0),$$
(1.8)

$$g^{(\text{xiv})}(0) = -135135f^{(\text{vii})}(z_0) - 945945f^{(\text{vi})}(z_0) - 1891890f^{(\text{v})}(z_0) - 1201200f^{(\text{iv})}(z_0) - 195195f^{'''}(z_0) - 4095f^{''}(z_0) - f^{\prime}(z_0),$$
(1.9)

etc.; the derivatives of odd order at the origin being at zero, since g(z) is an even function of z. Now, consider the expression

$$g(z) \approx A_1 \cos z - A_2 g(z) \cos z + A_3 \cos 2z - A_4 g(z) \cos 2z + \cdots + A_{2s-1} \cos sz - A_{2s} g(z) \cos sz, \qquad (1.10)$$

where the A_k 's are constants to be determined, and the \approx in (1.10) is to be interpreted in the sense that the Maclaurin expansions of both sides agree through the first 2s terms.

Note that both sides of (1.10) are, of course, even, as they should be.

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Observe that the Cauchy product of g(z) and $\cos mz$ is

$$g(z)\cos mz = \sum_{n=0}^{\infty} \sum_{k=0}^{n} \frac{g^{(2n-2k)}(0)(-1)^k m^{2k} z^{2n}}{(2n-2k)! (2k)!}.$$
 (1.11)

Since $\cos mz$ is entire, the above Cauchy product will have the same circle of convergence that equation (1.1) has (see [4]).

Using (1.11) to equate powers of z in (1.10) we find, after multiplying through by $(-1)^n (2n)!$,

$$(-1)^{n} g^{(2n)}(0) = A_{1} - A_{2} \sum_{k=0}^{n} (-1)^{n-k} {2n \choose 2k} g^{(2n-2k)}(0) + 2^{2n} A_{3}$$
$$- A_{4} \sum_{k=0}^{n} (-1)^{n-k} 2^{2k} {2n \choose 2k} g^{(2n-2k)}(0) + \dots + s^{2n} A_{2s-1}$$
$$- A_{2s} \sum_{k=0}^{n} (-1)^{n-k} s^{2k} {2n \choose 2k} g^{(2n-2k)}(0), \qquad (1.12)$$

where $\binom{n}{k}$ is the binomial coefficient.

Letting n = 0, 1, 2, ..., 2s - 1 in (1.12), we find an algebraic system of 2s equations with 2s unknowns for the determination of the A's. Then, g(z) is found as

$$g(z) \approx \frac{A_1 \cos z + A_3 \cos 2z + \dots + A_{2s-1} \cos sz}{1 + A_2 \cos z + A_4 \cos 2z + \dots + A_{2s} \cos sz}.$$
 (1.13)

Now, in equation (1.13), replace the above z by $\cos^{-1}(z - z_0 + 1)$, and make use of the defining equation for Chebyshev polynomials of the first kind $T_n(z) = \cos(n \cos^{-1} z)$, recalling the relation between f(z) and g(z) to obtain

$$f(z) \approx \frac{A_1 T_1 (z - z_0 + 1) + A_3 T_2 (z - z_0 + 1) + \dots + A_{2s-1} T_s (z - z_0 + 1)}{T_0 (z - z_0 + 1) + A_2 T_1 (z - z_0 + 1) + \dots + A_{2s} T_s (z - z_0 + 1)},$$
(1.14)

which gives a rational Chebyshev approximation of f(z) where the only restriction which has been assumed is analyticity of the function at z_0 .

Power series of the form given in (1.1) are sometimes found Taylor-made in the literature. For instance, see [6],

$$\exp(\cos z - 1) = 1 - \frac{1}{2}z^2 + \frac{1}{6}z^4 - \frac{31}{720}z^6 + \cdots, \qquad (1.15)$$

where the general coefficient is

$$\frac{(-1)^{n} 2^{1-n}}{n! (2n)!} \sum_{k=0}^{n-1} (-1)^{k} 2^{k} (-n)_{k} \sum_{r=0}^{n-k-1} \frac{(2k-2n)_{r}}{r!} (n-k-r)_{2n}, \qquad (1.16)$$

where $(\alpha)_n = \alpha(\alpha+1)(\alpha+2)\cdots(\alpha+n-1)$, $(\alpha)_0 = 1$, $\alpha \neq 0$, is Pochhammer's symbol. In series (1.15), $z_0 = 0$.

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Also, see [5],

$$\log \cos z = \sum_{n=1}^{\infty} (-1)^n (2^{2n} - 1) 2^{2n-1} B_{2n} z^{2n} / [n(2n)!], \qquad (1.17)$$

where the B_{2n} are Bernoulli numbers (see [1]). In the series (1.17), $z_0 = 1$.

It will be noticed that the coefficient of $f^{(j)}(z_0)$ in the sum for $g^{(2i)}(0)$, (i = 1, 2, ..., 2s - 1, j = 1, 2, ..., 2s - 1), exemplified in the list given at the beginning of this section, equations (1.2) through (1.9), is also the coefficient of $\cos jz$, evaluated at z = 0, in

$$\frac{d^{2j}}{dz^{2j}} \left(\exp(\cos z - 1) \right)$$

This provides a simple computer algorithm for generating these coefficients. This observation is due to one of the authors (Rosenthal).

2. ADAPTING THE ALGORITHM FOR THE GENERALIZED HYPERGEOMETRIC FUNCTION

The method we have developed enables us to find, in simple fashion, a rational Chebyshev approximation for the generalized hypergeometric function ${}_{p}F_{q}(z)$:

$${}_{P}F_{q}(a_{1}, a_{2}, \dots, a_{p}; b_{1}, b_{2}, \dots, b_{q}; z) = 1 + \sum_{n=1}^{\infty} \frac{(a_{1})_{n}(a_{2})_{n} \cdots (a_{p})_{n} z^{n}}{(b_{1})_{n}(b_{2})_{n} \cdots (b_{q})_{n} n!},$$
(2.1)

where none of the b's is zero or a negative integer (see [14]).

The derivative of (2.1) is given by (see [14])

$$\frac{a_1 a_2 \cdots a_p}{b_1 b_2 \cdots b_q} {}_p F_q(a_1 + 1, a_2 + 1, \dots, a_p + 1; b_1 + 1, b_2 + 1, \dots, b_q + 1; z).$$
(2.2)

The value of the hypergeometric function at the origin is 1. Hence, choosing $z_0 = 0$, it is quite simple to determine successive derivatives of the ${}_pF_q(z)$ at the origin to find, with the aid of equations (1.2) through (1.9), the values of g(z) and its derivatives at z = 0.

Note that g(0) and its derivatives at the origin will be given as rational functions of the coefficients of the ${}_{p}F_{q}(z)$. In particular, if these coefficients are themselves rational, then the rational Chebyshev approximation will involve only rational coefficients.

As the reader no doubt knows, many known functions are special cases (at most with a multiplicative monomial) of the generalized hypergeometric function. We will choose Bessel functions,

$$J_n(z) = \frac{(z/2)^n}{\Gamma(1+n)} {}_0F_1\left(-; 1+n; -\frac{1}{4}z^2\right),$$
(2.3)

to illustrate the algorithm.

It will be recalled that we mentioned, following (2.2), that, if the parameters appearing in the hypergeometric function are rational numbers, then the A's, the solutions of the system of equations (1.12), are also rational numbers. This holds true in most of the important cases. For this reason, we found it desirable to make use of a program (we chose REDUCE [15]) that did not execute the operation of division, so that the A's would be given in fractional form.

We close this section by making a comment that is probably obvious to the reader. If one wishes to go from a given s, the highest order of the Chebyshev polynomials in (1.14), to s+1 in the system of equations (1.12), then the matrix of the coefficients for s+1 will be the same as that for s, except that two rows and two columns will be added. Hence, knowing the inverse of the $2s \times 2s$ matrix one can find the inverse of the $(2s+2) \times (2s+2)$ matrix by using the method of partitioning in the technique known as "inversion by bordering."

3. ILLUSTRATING THE ALGORITHM

We will now give some examples of rational Chebyshev approximations obtained by use of the procedure outlined in the previous section. To list the approximations, we will give them in the following format:

$$f(z) \approx \frac{az^{k} (p_{0} z^{n} + p_{1} z^{n-1} + p_{2} z^{n-2} + \dots + p_{n-1} z + p_{n})}{b(q_{0} z^{m} + q_{1} z^{m-1} + q_{2} z^{m-2} + \dots + q_{m-1} z + q_{m})},$$
(3.1)

where $k + n \le s$, and $m \le s$. For each s we will simply list the coefficients in (3.1).

$$\frac{f(z) = J_0(z)}{s = 2}$$

 $a = 4, k = 0, n = 2, p_0 = 2, p_1 = 0, p_2 = -3;$
 $b = 1, m = 2, q_0 = 5, q_1 = 0, q_2 = -12.$

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 $a = -1, k = 0, n = 3, p_0 = 69, p_1 = 51, p_2 = -368, p_3 = -272;$ $b = 1, m = 3, q_0 = 23, q_1 = 17, q_2 = 368, q_3 = 272.$

s = 3

 $\underline{s=6}$

$$a = 12, k = 0, n = 6, p_0 = 57\ 76742, p_1 = 0, p_2 = -1838\ 79735,$$

 $p_3 = 0, p_4 = 10070\ 89152, p_5 = 0, p_6 = -7895\ 61600;$

$$b = 1, m = 6, q_0 = 60\,35647, q_1 = 0, q_2 = 3705\,82236, q_3 = 0,$$

 $q_4 = 97163\,85024, q_5 = 0, q_6 = -94747\,39200.$

The reader should observe that the magnitude of the coefficients increases quite rapidly with increasing s. We shall shortly see that the quality of the approximation also improves very rapidly as s increases.

$$s = 10$$

$$a = 300, k = 0, n = 10$$

$$\begin{array}{ll} p_0 = 2114\ 63570\ 00545\ 36614, & p_1 = 0, \\ p_2 = -4\ 28033\ 75450\ 19518\ 86781, & p_3 = 0, \\ p_4 = 28117868\ 03665\ 81890\ 18624, & p_5 = 0, \\ p_6 = -6194\ 13498\ 85928\ 62663\ 77984, & p_7 = 0, \\ p_8 = 31326\ 83622\ 73236\ 69829\ 38624, & p_9 = 0, \\ p_{10} = -23739\ 05902\ 96182\ 29215\ 88736; & \end{array}$$

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$$b = 1, m = 10,$$

$q_0 = 3272566141496807057,$	$q_1 = 0$,
$q_2 = 984654951486417966500,$	$q_3 = 0$,
$q_4 = 1597\ 67150\ 04330\ 42594\ 24000,$	$q_5 = 0$,
$q_6 = 15744170286741008972160000,$	$q_7 = 0$,
$q_8 = 761762144097573375762432000,$	$q_{9} = 0,$
$q_{10} = -712171770888546876476620800.$	

4. NUMERICAL VALUES AND GRAPHS OF SOME RATIONAL CHEBYSHEV APPROXIMATIONS

In this section we present the results of evaluating the rational forms given in section 3. The runs for different values of the parameter s will be contrasted with the tabulated values given in [1]. The latter will be taken, for purposes of comparison, as exact values.

 $f(z) = J_0(z)$

			•	., .							
Exact values			s = 2					<i>s</i> = 3			
Z	$J_0(z)$		Z		$J_0(z)$		Z		$J_0(z)$		
0.0	1.00000 00000	00000	0.0	1.00000	00000	00000	0.0	1.00000	00000	00000	
0.1	0.99750 15620	66040	0.1	0.99748	95397	48954	0.1	0.99750	15615	24048	
0.2	0.99002 49722	39576	0.2	0.98983	05084	74576	0.2	0.99002	49376	55860	
0.3	0.97762 62465	38296	0.3	0.97662	33766	23377	0.3	0.97762	58545	68055	
0.4	0.96039 82266	59563	0.4	0.95714	28571	42857	0.4	0.96039	60396	03960	
0.5	0.93846 98072	40813	0.5	0.93023	25581	39535	0.5	0.93846	15384	61539	
0.6	0.91200 48634	97211	0.6	0.89411	76470	58824	0.6	0.91198	04400	97800	
0.7	0.88120 08886	07405	0.7	0.84607	32984	29319	0.7	0.88114	00848	99939	
0.8	0.84628 73527	50480	0.8	0.78181	81818	18182	0.8	0.84615	38461	53846	
0.9	0.80752 37981	22545	0.9	0.69433	96226	41509	0.9	0.80725	75847	70970	
1.0	0.76519 76865	57967	1.0	0.57142	85714	28571	1.0	0.76470	58823	52941	
1.1	0.71962 20185	27511	1.1	0.38991	59663	86554	1.1	0.71876	81580	47647	
1.2	0.67113 27442	64363	1.2	0.09999	99999	99999	1.2	0.66972	47706	42202	
1.3	0.62008 59895	61509	1.3	-0.42816	90140	84507	1.3	0.61786	31995	47767	
1.4	0.56685 51203	74289	1.4	-1.67272	72727	27273	1.4	0.56347	43875	27840	
1.5	0.51182 76717	35918	1.5	-8.00000	00000	00000	1.5	0.50684	93150	68493	
1.6	0.45540 21676	39381	1.6	10.60000	00000	00000	1.6	0.44827	58620	68966	
1.7	0.39798 48594	46109	1.7	4.53877	55102	04082	1.7	0.38803	59978	82478	
1.8	0.33998 64110	42558	1.8	3.31428	57142	85714	1.8	0.32640	33264	03326	
1.9	0.28181 85593	74385	1.9	2.79008	26446	28099	1.9	0.26364	09994	90056	
2.0	0.22389 07791	41236	2.0	2.50000	00000	00000	2.0	0.20000	00000	00000	
2.1	0.16660 69803	31990	2.1	2.31641	79104	47761	2.1	0.13571	77853	99314	
2.2	0.11036 22669	22174	2.2	2.19016	39344	26230	2.2	0.07101	72744	72169	
2.3	0.05553 97844	45602	2.3	2.09826	98961	93771	2.3	0.00610	61531	23532	
2.4	0.00250 76832	97244	2.4	2.02857	14285	71429	2.4	-0.05882	35294	11765	
2.5	-0.04838 37764	68198	2.5	1.97402	59740	25974	2.5	-0.12359	55056	17978	
			Error	at $z = 0$.1: 1.2	20E-05	Error	at $z = 0$	1: 5.4	2E-10	
			Error	at $z = 1$	5: 8.	512	Error	at $z = 1$.5: 0.0	005	
			Error	at $z = 2$	5:-2.0	022	Error	at $z = 2$.5: 0.0	075	

	s =	6			<i>s</i> =	10	
Z	J	$f_0(z)$		Z		$J_0(z)$	
0.0	1.00000	00000	00000	0.0	1.00000	00000	00000
0.1	0.99750	15620	66040	0.1	0.99750	15620	66040
0.2	0.99002	49722	39576	0.2	0.99002	49722	39576
0.3	0.97762	62465	38249	0.3	0.97762	62465	38296
0.4	0.96039	82266	57938	0.4	0.96039	82266	59564
0.5	0.93846	98072	14225	0.5	0.93846	98072	40813
0.6	0.91200	48632	17224	0.6	0.91200	48634	97211
0.7	0.88120	08863	32675	0.7	0.88120	08886	07405
0.8	0.84628	73359	87120	0.8	0.84628	73527	50480
0.9	0.80752	36414	25774	0.9	0.80752	37981	22545
1.0	0.76519	88991	81058	1.0	0.76519	76865	57967
1.1	0.71962	28437	77746	1.1	0.71962	20185	27512
1.2	0.67113	39900	87712	1.2	0.67113	27442	64364
1.3	0.62008	81243	85951	1.3	0.62008	59895	61514
1.4	0.56685	88740	92599	1.4	0.56685	51203	74305
1.5	0.51183	42373	87263	1.5	0.51182	76717	35967
1.6	0.45541	34601	77972	1.6	0.45540	21676	39523
1.7	0.39800	38749	36571	1.7	0.39798	48594	46502
1.8	0.34001	77192	20127	1.8	0.33998	64110	43589
1.9	0.28186	89650	63377	1.9	0.28181	85593	76972
2.0	0.22397	01919	55021	2.0	0.22389	07791	47447
2.1	0.16672	95358	25093	2.1	0.16660	69803	46316
2.2	0.11054	77454	23837	2.2	0.11036	22669	54003
2.3	0.05581	53758	52507	2.3	0.05553	97845	13916
2.4	0.00291	01468	75270	2.4	0.00250	76834	39234
2.5	-0.04780	55089	54713	2.5	-0.04838	37761	81732
Error	at $z = 0$	1: 0		Error	at $z = 0$	1: 0	
Error	at $z = 1$	5:-6.5	57E-06	Error	at $z = 1$	5:-4.9	90E-14
Error	at $z = 2$	5:-5.7	78E-04	Error	at $z = 2$	5:-2.8	86E-10

The algorithm is seen to be very stable. As the value of s increases, the quality of the approximations improves notably. The last example above, $J_0(z)$ for s = 10, gives remarkable agreement throughout the range $0 \le |z| \le 2.5$.

5. ZEROS OF THE DENOMINATOR POLYNOMIALS OF THE RATIONAL CHEBYSHEV APPROXIMATIONS

If in equation (1.10) we let s increase without bound, then both sides will represent the same function since their Maclaurin expansions agree for all terms. In this case, equation (1.13) will have an infinite series in both the numerator and denominator. The values of z for which the series in the denominator converges to zero will be singular points of g(z), unless the series in the numerator also converges to zero there. As equation (1.13) stands, it being an approximate relation, it is conceivable that the right-hand side may have poles which are not singular points of the function g(z). This implies, of course, that the right-hand side of equation (1.14) may also have poles which are not singular points of f(z). These would be the so-called *spurious poles*. Let us look at this phenomenon somewhat more closely for the example given in Section 3.

The denominator polynomial of the rational Chebyshev approximation for the Bessel function $J_0(z)$ corresponding to s = 10 has real zeros at the points

$z = \pm 0.95778$ 12766 24968 22726 05909 45945.

Yet, the graph given in Figure 1, and the table of values of this function do not seem to indicate any abnormal behavior in the neighborhood of this point. However, if we analyze the rational approximation within $\pm E$ -18 of this point, then the rational form is seen to undergo marked oscil-

lations with nearly infinite slope. Nevertheless, as soon as we are within $\pm E-17$ of the point in question, the erratic behavior disappears and the algorithm again represents the correct values of the Bessel function $J_0(z)$.





This figure shows the Bessel function of the first kind of order zero, $J_0(z)$ plotted against the rational Chebyshev approximation corresponding to s = 10. After z = 9, the Bessel function continues to oscillate, while the approximation separates from this behavior. The two functions move apart after z = 7. The algorithm approximates the first zero of the Bessel function to be 2.40482 55580, and the second zero to be 5.51960 87207. These results compare favorably with the correct values 2.40482 55577 and 5.52007 81103 given in [1].

We shall now speak of the significance of these roots. The highly localized character of the oscillation indicates that the numerator polynomial also has zeros which are very close to the zeros of the denominator polynomial. This is indeed the case for all of the examples we studied. The numerator polynomial of the s = 10 approximation of the Bessel function, for instance, has real zeros at the points

$z = \pm 0.95778$ 12766 24968 22150 32913 84229

which match the zeros of the denominator polynomial through seventeen decimal places. The oscillatory behavior is then simply a reflection of the computer's arithmetic inability to handle 0/0. The algorithm, we see, is a self-correcting one that introduces zeros in the numerator and denominator polynomials in a way that ensures the correct approximation to the function for a given value of *s*.

In essence, our method provides a rational approximation $P_s(z)/Q_s(z)$ such that its Taylor expansion about the point z_0 agrees with the Taylor expansion of f(z) through the first 2s terms. This requirement may be written as

$$Q_s(z)f(z) - P_s(z) = (z - z_0)^{2s+1} \sum_{k=0}^{\infty} c_k (z - z_0)^k$$

and it is equivalent to the criterion for choosing the s^{th} diagonal entry in the Padé table for $z_0 = 0$.

Because of the proximity of the real zeros of the numerator and denominator polynomials of the Bessel function approximation corresponding to s = 10, we chose to divide out the zeros and try out the outcome against the tabulated values given before. The resulting expression is:

$$a = 300, k = 0, n = 8$$

$p_0 = 2114\ 63570\ 00545\ 36614\ 00000\ 00000\ 0,$	$p_1 = 0,$
$p_2 = -426093904070975989175.46176795488,$	$p_3 = 0,$
$p_4 = 277\ 26992\ 93536\ 91260\ 65928\ 27801\ 88600\ 0,$	$p_{5} = 0,$
$p_6 = -593978281249957947189316.44693046000,$	$p_7 = 0$,
$p_8 = 25878\ 00631\ 84966\ 42221\ 73861\ 61878\ 68000\ 0,$., ,

b = 1, m = 8

$q_0 = 3272\ 56614\ 14968\ 07057\ 00000\ 00000\ 0,$	$q_1 = 0$,
$q_2 = 987657023587922726257.68351763259,$	$q_3 = 0$,
$q_4 = 1606\ 73172\ 24978\ 36037\ 41999.24516\ 09500\ 0,$	$q_{5} = 0$,
$q_6 = 15891563013737422140956470.02415100000,$	$q_7 = 0$,
$q_8 = 776340189554899257318873022.34814000000$	1, ,

The tabulated values resulting from this approximation are:

Z	$J_0(z)$
0.0	1.00000 00000 00000
0.1	0.99750 15620 66040
0.2	0.99002 49722 39576
0.3	0.97762 62465 38296
0.4	0.96039 82266 59564
0.5	0.93846 98072 40813
0.6	0.91200 48634 97211
0.7	0.88120 08886 07405
0.8	0.84628 73527 50480
0.9	0.80752 37981 22545
1.0	0.76519 76865 57967
1.1	0.71962 20185 27512
1.2	0.67113 27442 64364
1.3	0.62008 59895 61514
1.4	0.56685 51203 74305
1.5	0.51182 76717 35967
1.6	0.45540 21676 39523
1.7	0.39798 48594 46502
1.8	0.33998 64110 43589
1.9	0.28181 85593 76972
2.0	0.22389 07791 47447
2.1	0.16660 69803 46316
2.2	0.11036 22669 54003
2.3	0.05553 97845 13916
2.4	0.00250 76834 39234
2.5	-0.04838 37761 81732
Error	at $z = 0.1: 0$
Error	at $z = 1.5:-4.90E-14$
Error	at $z = 2.5:-2.86E-10$

These are exactly the same values, to fifteen-decimal accuracy, obtained with the s = 10 approximation of the Bessel function $J_0(z)$ before the roots are divided out!—These results imply a substantial saving in computer time since the number of divisions required for a given approximation is reduced by two.

A comment is in order, though it is probably obvious to the reader. The results shown in the above table were obtained by dividing the numerator polynomial by its real roots, and the denominator polynomial by its corresponding real roots. Slightly better accuracy is obtained (though the

above table does not indicate it) if we divide both numerator and denominator polynomials by either the real roots of the numerator *or* the real roots of the denominator since, in this case, all we are doing is dividing numerator and denominator of the s = 10 approximation by a common factor.

It is worth emphasizing that the rational Chebyshev approximations our algorithm provides are not optimal, in the sense that error does not remain constant within the range of approximation. Rather, error is least when one is sufficiently near the point z_0 and the quality of the approximation deteriorates as we move away from the point in question. The importance of the method lies, we believe, in the extreme simplicity with which it can provide rational Chebyshev approximations of any accuracy for a wide variety of functions. These nonoptimal approximations may easily be used to obtain optimal Chebyshev approximations. Several algorithms have been developed to this effect.

Let us speak now of the origin of the problem that has occupied us in the last five sections.

6. SOME HISTORY

About one hundred and twenty-five years ago, the Russian mathematician Pafnuty Lvovich Chebyshev (1821-1894) set himself the problem of finding the best rational approximation of a continuous function specified on an interval [a, b]. Specifically, he wanted to determine parameters $p_0, p_1, ..., p_n$; $q_0, q_1, ..., q_m$ in the expression

$$Q(x) = s(x) \frac{p_0 x^n + p_1 x^{n-1} + \dots + p_n}{q_0 x^m + q_1 x^{m-1} + \dots + q_m},$$
(6.1)

where *m* and *n* are given, and s(x) is a function continuous on [a, b], so that the deviation of Q(x) from a chosen continuous function f(x),

$$H_{\mathcal{Q}} = \max_{\substack{a \le x \le b}} |f(x) - Q(x)| \tag{6.2}$$

shall be a minimum.

Chebyshev established the beautiful existence theorem [6; 2]:

The function P(x), which deviates least from the function f(x) than does any other function of the type exemplified by equation (6.1) is completely characterized by the following property: If the function can be expressed in the form

$$P(x) = s(x)\frac{a_0 x^{n-\sigma} + a_1 x^{n-\sigma-1} + \dots + a_{n-\sigma}}{b_0 x^{m-\tau} + b_1 x^{m-\tau-1} + \dots + b_{m-\tau}} = s(x)\frac{A(x)}{B(x)}$$

where $0 \le \sigma \le n$, $0 \le \tau \le m$, $b_0 \ne 0$ and the fraction $\frac{A(x)}{B(x)}$ is irreducible, then the number N of consecutive points of the interval [a, b] at which the difference f(x) - P(x), with alternate change of sign, takes on the value H_p , is not less than m+n+2-d, where $d = \min(\sigma, \tau)$; in case $P(x) \equiv 0$, then $N \ge n+2$.

Chebyshev did not provide a constructive approach to the problem of finding the rational approximations whose existence is guaranteed by the above theorem. He, and E. Solotarev did work out one example, based on the theory of Jacobian elliptic functions, that meets the require-

ments of the theorem [16]. Since that time, though, many people have sought to obtain an explicit method of attack for determining these rational approximations [8; 9; 10]. The problem is especially complicated by the fact that the class of continuous functions is a very broad one. Most of the methods of attack that have been developed deal with a more restrictive class of functions: bounded variation, analytic, or the like.

A substantial advance was made by H. Padé in his now classic thesis of 1892 [13]. Padé's method, mentioned briefly at the end of the last section, yields excellent rational approximations of analytic functions by means of solutions of a system of linear algebraic equations [18]. The method is an extension of some earlier work of Frobenius [6]. However, it does not provide rational Chebyshev approximations. It is known that rational forms in Chebyshev polynomials yield better accuracy than ordinary rational forms [16].

Maehly gave a method for obtaining rational Chebyshev approximations of functions of bounded variation on the unit interval [12; 16]. It has the substantial disadvantage of requiring that the given function be first expanded in a series of Chebyshev polynomials. If the function is anywhere complicated, these expansions may be devilishly hard to obtain.

To the best of our knowledge, no method is known for obtaining rational Chebyshev approximations that is better, more direct, or more powerful than the one we have presented in this paper. The method was discovered by one of the authors (Castellanos) as a result of his work on formulas to approximate π while in preparation of "The Ubiquitous π ," *Math. Magazine* **61.2-3** (April-June 1988). The delicate and time-consuming task of carrying the algorithm into a working computer program was done by the other author (Rosenthal).

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