

A RELIABILITY PROBLEM

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ABSTRACT

An $m \times n$ array of elements is considered in which each element has a probability p of being reliable. The array as a whole is considered reliable if there does not exist in the array any polydominoe of a given form in any orientation having all of its elements unreliable. A method is given for determining the probability of reliability for the array and solutions are worked out explicitly for several special cases.

1. INTRODUCTION

We are given an $m \times n$ array

	1	2	...	n
1				
2				
⋮				
m				

in which each of the mn elements has a given probability p of being reliable (and a probability q of being unreliable where $p + q = 1$). The $m \times n$ matrix as a whole will be considered reliable, if and only if, there does not exist in the array, any polydominoe of a given form in any orientation having all of its elements unreliable. The problem then is to calculate the probability of reliability of the array. The special case where $m = 2$ and the given polydominoe is a 2×1 arose in the design of a low-altitude detection antenna. If any 2×1 polydominoe had both its elements unreliable, then the antenna could not fulfill its detection mission.

The specific cases to be considered here explicitly are the following:

	Array size	Failure polydominoe		
(C ₁)	$2 \times n$	<table border="1" style="display: inline-table; border-collapse: collapse;"> <tr> <td style="padding: 2px 5px;">q</td> <td style="padding: 2px 5px;">q</td> </tr> </table>	q	q
q	q			

	Array size	Failure polydominoe				
(C ₂)	2 × n	<table border="1" style="display: inline-table; vertical-align: middle;"><tr><td>q</td></tr><tr><td>q</td><td>q</td></tr></table>	q	q	q	
q						
q	q					
(C ₃)	2 × n	<table border="1" style="display: inline-table; vertical-align: middle;"><tr><td>q</td><td>q</td></tr><tr><td>q</td><td>q</td></tr></table>	q	q	q	q
q	q					
q	q					
(C ₄)	3 × n	<table border="1" style="display: inline-table; vertical-align: middle;"><tr><td>q</td><td>q</td></tr></table>	q	q		
q	q					
(C ₅)	3 × n	<table border="1" style="display: inline-table; vertical-align: middle;"><tr><td>q</td><td>q</td></tr><tr><td>q</td><td>q</td></tr></table>	q	q	q	q
q	q					
q	q					

2. PROBABILITY FOR RELIABILITY

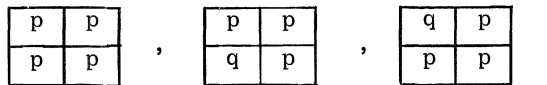
For all the cases, we will let P_n denote the probability of the 2 × n or 3 × n array being reliable. For the 2 × n array, A_n, B_n, C_n, D_n will denote the respective probabilities of reliability of the array if the end 2 × 1 polydominoe has the form



and then

$$(1) \quad P_n = A_n + B_n + C_n + D_n .$$

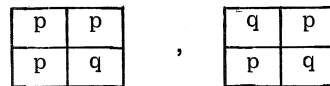
For the case (C₁), D_n = 0 and B_n = C_n. Here, for an A_{n+1} array, the end 2 × 2 polydominoe must have one of the three following forms:



Thus,

$$(2) \quad A_{n+1} = p^2(A_n + B_n + C_n) .$$

For a B_{n+1} array, the end 2 × 2 polydominoe must have one of the two forms



and thus

$$(3) \quad B_{n+1} = pq(A_n + C_n)$$

and similarly

$$(4) \quad C_{n+1} = pq(A_n + B_n) .$$

On eliminating B_n and C_n, we obtain

$$(5) \quad A_{n+1} = p^2 P_n ,$$

$$(6) \quad P_{n+1} = pP_n + pqA_n ,$$

and then

$$(7) \quad P_{n+1} = pP_n + p^3 q P_{n-1} .$$

For initial conditions, we have

$$(8) \quad A_1 = p^2, \quad B_1 = pq = C_1, \quad D_1 = 0.$$

Whence,

$$(9) \quad P_1 = 1 - q^2, \quad P_2 = 2p^2 - p^4 .$$

The solution of (7) is then given by

$$P_n = k_1 r_1^n + k_2 r_2^n ,$$

where r_1, r_2 are the roots of $x^2 = px + p^3q$ and constants k_1, k_2 are determined so as to satisfy (9). This gives

$$(10) \quad P_n = \frac{1 - q^2}{a} \left\{ \left(\frac{p+a}{2} \right)^n - \left(\frac{p-a}{2} \right)^n \right\} + \frac{p^3q}{a} \left\{ \left(\frac{p+a}{2} \right)^{n-1} - \left(\frac{p-a}{2} \right)^{n-1} \right\} ,$$

where $a = \sqrt{p^2 + 4p^3q}$.

For (C₂), it then follows as before that

$$(11) \quad A_{n+1} = p^2(A_n + B_n + C_n + D_n) ,$$

$$(12) \quad B_{n+1} = C_{n+1} = pq(A_n + B_n + C_n) ,$$

$$(13) \quad D_{n+1} = q^2 A_n$$

subject to the initial conditions,

$$(14) \quad A_1 = p^2, \quad B_1 = pq = C_1, \quad D_1 = q^2 .$$

Eliminating B_n, C_n, D_n , we obtain

$$(15) \quad A_{n+2} = (p^2 + 2pq)A_{n+1} + p^2q^2A_n - 2p^3q^3A_{n-1} .$$

Whence,

$$A_n = k_1 r_1^n + k_2 r_2^n + k_3 r_3^n ,$$

where r_1, r_2, r_3 are the roots of

$$x^3 = (p^2 + 2pq)x^2 + p^2q^2x - 2p^3q^3$$

and the constants k_1, k_2, k_3 are determined from the initial conditions (note that here $A_1 = A_2 = p^2, A_3 = p^4[1 + 2pq + 5q^2]$). Then B_n, C_n, D_n and P_n are easily determined.

For (C₃), we have (11) and

$$(16) \quad B_{n+1} = C_{n+1} = pq(A_n + B_n + C_n + D_n) ,$$

$$(17) \quad D_{n+1} = q^2(A_n + B_n + C_n)$$

(again all subject to conditions (14)). On eliminating D_n , we obtain

$$(18) \quad A_{n+1} = p^2 A_n + B_n + C_n + q^2(A_{n-1} + B_{n-1} + C_{n-1}) ,$$

$$(19) \quad pB_{n+1} = pC_{n+1} = qA_{n+1} .$$

Whence,

$$(20) \quad A_{n+1} = p(p + 2q)(A_n + q^2A_{n-1}) .$$

Then,

$$A_n = k_1 r_1^n + k_2 r_2^n ,$$

where r_i are the roots of

$$x^2 = p(p + 2q)(x + q)$$

and the k_i 's are determined from the initial conditions.

Then B_n , C_n , D_n and P_n are found from (14), (17) and (1).

For the $3 \times n$ arrays, we let A_n , B_n , C_n , D_n , E_n , F_n , G_n , H_n denote the respective probabilities of reliability of the array if the end 3×1 polydominoe has the form

$$\begin{array}{|c|} \hline p \\ \hline p \\ \hline p \\ \hline \end{array} , \begin{array}{|c|} \hline p \\ \hline p \\ \hline q \\ \hline \end{array} , \begin{array}{|c|} \hline p \\ \hline q \\ \hline p \\ \hline \end{array} , \begin{array}{|c|} \hline q \\ \hline p \\ \hline p \\ \hline \end{array} , \begin{array}{|c|} \hline q \\ \hline q \\ \hline p \\ \hline \end{array} , \begin{array}{|c|} \hline q \\ \hline p \\ \hline q \\ \hline \end{array} , \begin{array}{|c|} \hline p \\ \hline q \\ \hline q \\ \hline \end{array} , \begin{array}{|c|} \hline q \\ \hline q \\ \hline q \\ \hline \end{array}$$

and

$$(21) \quad P_n = A_n + B_n + C_n + D_n + E_n + F_n + G_n + H_n .$$

For (C_4) ,

$$E_n = G_n = H_n = 0, \quad B_n = D_n ,$$

$$(22) \quad A_1 = p^3, \quad B_1 = C_1 = D_1 = p^2q, \quad F_1 = pq^2 ,$$

$$(23) \quad A_{n+1} = p^3(A_n + B_n + C_n + D_n + F_n) ,$$

$$(24) \quad B_{n+1} = D_{n+1} = p^2q(A_n + C_n + D_n) ,$$

$$(25) \quad C_{n+1} = p^2q(A_n + B_n + D_n + F_n) ,$$

$$(26) \quad F_{n+1} = pq^2(A_n + C_n) .$$

For (C_5) ,

$$B_n = C_n = D_n, \quad E_n = G_n,$$

$$(27) \quad A_1 = p^3, \quad B_1 = C_1 = D_1 = p^2q, \quad E_1 = F_1 = G_1 = pq^2, \quad H_1 = q^3 ,$$

$$(28) \quad A_{n+1} = p^3 P_n ,$$

$$(29) \quad B_{n+1} = p^2 q P_n ,$$

$$(30) \quad E_{n+1} = pq^2(A_n + B_n + C_n + D_n + F_n + G_n) ,$$

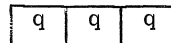
$$(31) \quad F_{n+1} = pq^2 P_n ,$$

$$(32) \quad H_{n+1} = q^3(A_n + B_n + C_n + D_n + F_n) .$$

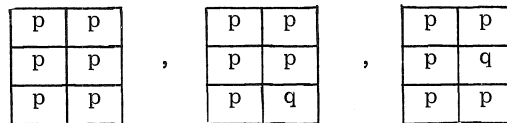
Although we can carry out the elimination process for (C_4) and (C_5) by means of the operator E and then determine P_n in terms of the roots of a higher order polynomial, it is not worthwhile. In these cases (and even some of the prior ones), one can just use a computer on the recurrence relations to determine the P_n 's.

3. HIGHER ORDER POLYDOMINOES

The previous methods, with some adaptation, will also apply when the failure polydominoe is of higher order than the previous ones. As in the last two cases, it will suffice to just get the appropriate recurrence equations. If the failure polydominoe is of the type



in a $3 \times n$ array, then we would need terms A_n, B_n, \dots , corresponding to a reliable $3 \times n$ array whose end 2×3 polydominoe has the forms



etc.

This will, of course, lead to an increased number of recurrence relations. Other arrays which can be solved similarly are cylindrical and torodial ones as well as higher dimensional ones.

