Modal and tense operators on a complete Boolean algebra

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Special Session on Algebraic Logic Western Sectional Meeting August 24, 2025 Definition. A **Stone space** X is a compact Hausdorff space such that every connected component is a point.

Definition. A **continuous relation** on X is a closed subset $R \subseteq X \times X$ such that

- 1. R[x] is closed for all points $x \in X$,
- 2. $R^{-1}[U]$ is clopen for all clopens $U \subseteq X$.

(Continuous relations on X generalize continuous functions on X.)

Definition.

$$\mathcal{CR}(X) = \{R \mid R \text{ is a continuous relation on } X\}$$

$$R \leq S \qquad \Longleftrightarrow \qquad R \subseteq S$$

Proposition. CR(X) has joins.

$$R \lor S = R \cup S$$

Let $\alpha \mathbb{N}$ be the one-point compactification of \mathbb{N} .

Proposition. $CR(\alpha \mathbb{N})$ does not have meets.

Proposition. $CR(\alpha \mathbb{N})$ is not distributive as a join-semilattice.

Definition. A join-semilattice *L* is **distributive** if

$$a \le b \lor c \implies (\exists b' \le b) (\exists c' \le c) \ a = b' \lor c'.$$

Definition. A Stone space X is **extremally disconnected** if the closure of every open set is open.

Proposition. In this case, CR(X) is a complete lattice.

$$\bigvee_{i\in I}R_i=\operatorname{cl}\left(\bigcup_{i\in I}R_i\right)$$

Theorem(Bezhanishvili, K). In this case, CR(X) is a coframe.

Definition. A complete lattice *L* is a **coframe** if

$$a \vee \left(\bigwedge_{i \in I} b_{\alpha}\right) = \bigwedge_{i \in I} a \vee b_{\alpha}.$$

Proposition. For general extremally disconnected X, the frame $\mathcal{CR}(X)^{op}$

- 1. is zero-dimensional,
- 2. need not be algebraic,
- 3. need not be spatial.

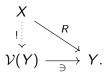
Let $\beta \mathbb{N}$ be the Stone-Čech compactification of \mathbb{N} .

Theorem(Bezhanishvili, K). $CR(\beta\mathbb{N}) \cong \Gamma(\mathbb{N} \times \beta\mathbb{N})$.

Corollary. The frame $\mathcal{CR}(\beta\mathbb{N})^{op}$ is

- 1. algebraic,
- spatial,
- 3. not a coframe,
- 4. not a Boolean algebra.

Definition. The **Vietoris space** V(Y) is $\Gamma(Y)$ with the topology such that



Theorem(Bezhanishvili, K). $CR(\beta \mathbb{N}) \cong \Gamma(\mathbb{N} \times \beta \mathbb{N})$.

Proof.

$$\mathcal{CR}(\beta\mathbb{N})\cong\mathcal{C}(\beta\mathbb{N},\mathcal{V}(\beta\mathbb{N}))\cong\mathcal{C}(\mathbb{N},\mathcal{V}(\beta\mathbb{N}))\cong\Gamma(\mathbb{N}\times\beta\mathbb{N})$$

Proposition. The function \cap : $\mathcal{V}(\beta\mathbb{N}) \times \mathcal{V}(\beta\mathbb{N}) \to \mathcal{V}(\beta\mathbb{N})$ is not continuous!

Definition. A **necessity operator** $\square \in \mathcal{NO}(B)$ on a Boolean algebra B is

$$\square \colon B \to B, \qquad \square 0 = 0, \qquad \square (a_1 \wedge a_2) = \square a_1 \wedge \square a_2.$$

Stone Duality. $B \cong \Omega(\operatorname{Spec}(B)) \cap \Gamma(\operatorname{Spec}(B))$.

Jónsson-Tarski Duality. $\mathcal{NO}(B) \cong \mathcal{CR}(\operatorname{Spec}(B))^{op}$.

$$\square_1 \leq \square_2 \qquad \Longleftrightarrow \qquad (\forall a \in B) \ \square_1 a \leq \square_2 a$$

Corollary. In general, $\mathcal{NO}(B)$ is a meet-semilattice that

- 1. need not have joins,
- 2. need not be distributive.

Corollary. When B is complete, $\mathcal{NO}(B)$ is a complete Heyting algebra.

Proposition. Let B be a complete Boolean algebra, and let $\Box \in \mathcal{NO}(B)$. The following are equivalent:

$$1. \ \Box\left(\bigwedge_{i\in I}a_i\right)=\bigwedge_{i\in I}\Box a_i,$$

2. $\Box a_1 \lor a_2 = 1$ iff $a_1 \lor \Box^{\dagger} a_2 = 1$ for some $\Box^{\dagger} \in \mathcal{NO}(B)$.

Definition. In this case, \square is a **tense necessity operator**, $\square \in \mathcal{TNO}(B)$.

Proposition. TNO(B) is a frame.

Theorem (Bezhanishvili, K). Let B be a complete Boolean algebra. The following are equivalent:

- 1. TNO(B) is a spatial frame,
- 2. TNO(B) is a complete atomic Boolean algebra,
- 3. B is atomic.

Definition. A complete Boolean algebra B is **measurable** if there exists

$$\mu \colon B \to [0, \infty]$$

- 1. $\mu(a) = 0$ iff a = 0,
- 2. $\mu(a_1 \lor a_2) = \mu(a_1) + \mu(a_2)$ whenever $a_1 \land a_2 = 0$,
- 3. $\mu\left(\bigvee_{n\in\omega}a_n\right)=\bigvee_{n\in\omega}\mu(a_n)$ whenever $a_0\leq a_1\leq\cdots$,
- 4. $0 < \mu(a_1)$ implies that $0 < \mu(a_0) < \infty$ for some $a_0 \le a_1$.

Example. The complete Boolean algebra $\mathcal{R}\Omega(\mathbb{R})$ is not measurable.

Theorem (folk?). The following are equivalent:

- B is measurable,
- ② $B \cong M(X)/N(X)$ for some decomposable measure space X.

Let B be a measurable complete Boolean algebra. We paraphrase.

Definition (Weaver). A **measurable relation** on B is a relation $R \subseteq B \times B$ such that

$$\left(\bigvee_{i\in I}a_i,\bigvee_{j\in J}b_j\right)\in R\iff (\exists i\in I)\ (\exists j\in J)\ (a_i,b_j)\in R.$$

Proposition. There is a one-to-one correspondence between

- 1. measurable relations on B,
- 2. tense necessity operators on B.

Theorem (Weaver). Let $\Box \in \mathcal{TNO}(B)$. Let $c_1, c_2 \in B$ with $\Box c_1 \lor c_2 \neq 1$. There exist complete Boolean algebra homomorphisms $\pi_i \colon B \to B'$ with

$$\pi_i(c_i) = 0,$$
 $\square a_1 \vee a_2 = 1 \implies \pi_1(a_1) \vee \pi_2(a_2) = 1.$

Let B be a measurable complete Boolean algebra. We paraphrase.

Definitions (Weaver). A tense necessity operator $\square \in \mathcal{TNO}(B)$ is

- 1. a measurable graph structure if $\square \leq id$ and $\square^{\dagger} = \square$,
- 2. a measurable poset structure if $\square \leq \square \circ \square$ and $\square \vee \square^{\dagger} = \mathrm{id}$,
- 3. a measurable equivalence relation if $\square \leq \mathrm{id}$, $\square^\dagger = \square$, and $\square^\dagger = \square$.
- 4. a measurable function if $id \leq \square \circ \square^{\dagger}$ and $\square^{\dagger} \circ \square \leq id$.

Proposition. There is a one-to-one correspondence between

- 1. measurable equivalence relations on B,
- 2. complete Boolean subalgebras of B.

Proposition. There is a one-to-one correspondence between

- 1. measurable functions on B,
- 2. complete Boolean algebra homomorphisms on *B*.

We compare two perspectives on complete Boolean algebras.

Spectral perspective.

- 1. Complete Boolean algebras are dual to extr. disc. Stone spaces.
- 2. Tense necessity operators are dual to interior relations.

Localic perspective.

- 1. Complete Boolean algebras are dual to discrete locales.
- 2. Tense necessity operators are dual to relations on those locales.

Complete Boolean algebras generalize sets in point-free topology.

Similarly, quantum OMLs generalize sets in noncommutative geometry.