

MATH 5330 ASSIGNMENT 6

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QUESTION 1

Solution: Let G be a k -regular graph of size n , so every vertex has degree k .

(a) We claim that $\chi_l(G) \leq k + 1$. To see this, assign lists of size $k + 1$ to each vertex of $V(G)$, let $V(G)$ be ordered with any order and colour greedily. When colouring v_i in this greedy colouring there will be at most k colours used in the neighbourhood of v_i and since v_i was assigned a list of size $k + 1$, there will be at least one colour remaining in v_i 's list that v_i can be coloured with. Therefore G is $(k+1)$ -choosable so $\chi_l(G) \leq k+1$.

(b) Suppose $L : V(G) \rightarrow \mathcal{P}A$ is a list assignment of G so that $|L(v)| = k$ for all $v \in V(G)$. Order $V(G)$ at random and apply the greedy algorithm. By this we mean for each v_i if there is a colour in its list that can be used to colour v_i , then colour it, otherwise leave it uncoloured. Define the random variable X to be the number of vertices that can be coloured in this way. For each $v \in V(G)$ define the indicator variable X_v as follows:

$$X_v = \begin{cases} 1 & \text{if } v \text{ can be coloured} \\ 0 & \text{otherwise} \end{cases}$$

Now $E(X_v) = \mathbb{P}(X_v = 1) \geq \mathbb{P}(X_v \text{ is ordered before at least one of its } k \text{ neighbours}) \geq \frac{k}{k+1}$. So the linearity of expectation gives:

$$\begin{aligned} E(X) &= \sum_{v \in V(G)} E(X_v) \\ &\geq \sum_{v \in V(G)} \frac{k}{k+1} \\ &= n \cdot \frac{k}{k+1} \end{aligned}$$

So by the pigeonhole property of expectation, there exists an ordering of $V(G)$ such that there exists a list colouring under L of at least $n \cdot \frac{k}{k+1}$ vertices of G .

(c) Suppose $1 \leq t \leq k$. Suppose $L : V(G) \rightarrow \mathcal{PA}$ is a list assignment of G so that $|L(v)| = t$ for all $v \in V(G)$. Order $V(G)$ at random and apply the greedy algorithm. Define the random variable X to be the number of vertices that can be coloured with the greedy algorithm. For each $v \in V(G)$ define the indicator variable X_v as follows:

$$X_v = \begin{cases} 1 & \text{if } v \text{ can be coloured} \\ 0 & \text{otherwise} \end{cases}$$

Now $E(X_v) = \mathbb{P}(X_v = 1) \geq \mathbb{P}(X_v \text{ is ordered before at least } k - t \text{ of its } k \text{ neighbours}) \geq \frac{t}{k+1}$. So the linearity of expectation gives:

$$\begin{aligned} E(X) &= \sum_{v \in V(G)} E(X_v) \\ &\geq \sum_{v \in V(G)} \frac{t}{k+1} \\ &= n \cdot \frac{t}{k+1} \end{aligned}$$

So by the pigeonhole property of expectation, there exists an ordering of $V(G)$ such that there exists a list colouring under L of at least $n \cdot \frac{t}{k+1}$ vertices of G .

QUESTION 2

Solution: Consider $G = G(n, p)$. Let X be the number of edges in G .

(a) For all subsets S such that $|S| = 2$ of $V(G)$ define the indicator variable X_S as follows:

$$X_S = \begin{cases} 1 & \text{if } G[S] \cong K_2 \\ 0 & \text{otherwise} \end{cases}$$

I.e, $X_S = 1$ if the vertices in S are joined by an edge. Now $X = \sum_{S \subseteq V(G), |S|=2} X_S$, so by the linearity of expectation,

$$E(X) = \sum_{S \subseteq V(G), |S|=2} E(X_S).$$

(a) We want to find $E(X_S) = \mathbb{P}(X_S = 1)$, which by definition of $G(n, p)$ is p . So from part (a), we have:

$$\begin{aligned} E(X) &= \sum_{S \subseteq V(G), |S|=2} E(X_S) \\ &= \sum_{S \subseteq V(G), |S|=2} p \\ &= \binom{n}{2} p \end{aligned}$$

(b) Markov's inequality gives:

$$\mathbb{P}(X \geq n) \leq \frac{E(X)}{n} = \frac{1}{n} \cdot \binom{n}{2} p$$

(c) Let p be a function of n , and assume that $p = o(\frac{1}{n})$. Now by (b), we know that

$$\begin{aligned} \mathbb{P}(X \geq n) &\leq \frac{1}{n} \cdot \binom{n}{2} p \\ &\leq np \\ &\rightarrow 0 \text{ as } n \rightarrow \infty \text{ (Since } p = o(\frac{1}{n}) \text{)} \end{aligned}$$

So a.a.s. $e(G) \leq n - 1$ and since a graph on n vertices must have at least $n - 1$ edges, a.a.s G is not connected.

(d) By definition of $G(n, p)$, the variables X'_S s are independent so $Cov(X_S, X_T) = 0$ if $S \neq T$ and therefore we have linearity of variance. Also, $Var(X_S) = E((X_S)^2) -$

$E(X_S)^2 = p - p^2$. Therefore,

$$\begin{aligned}
 \text{Var}(X) &= \text{Var}\left(\sum_{S \subseteq V(G), |S|=2} X_S\right) \\
 &= \sum_{S \subseteq V(G), |S|=2} \left(\text{Var} X_S\right) \\
 &= \sum_{S \subseteq V(G), |S|=2} (p - p^2) \\
 &= \binom{n}{2} p(1 - p) \\
 &= E(X)(1 - p)
 \end{aligned}$$

(e) We first note that $n - 1 = \frac{2E(X)}{np}$. Now,

$$\begin{aligned}
 \mathbb{P}(X \geq n) &= 1 - \mathbb{P}(X \leq n - 1) \\
 &\geq 1 - \mathbb{P}\left(|X - E(X)| \geq \frac{(2 - np)E(X)}{np}\right) \\
 &\geq 1 - \frac{(np)^2 \text{Var}(X)}{(np - 2)^2 E(X)^2} \text{ (by Chebeshev's inequality applied to } \mathbb{P}(X \leq n - 1)) \\
 &= 1 - \frac{n^2 p^2 (1 - p)}{(np - 2)^2 E(X)} \text{ (By part (d))}
 \end{aligned}$$

(f) Now, if p is constant, we see that:

$$\begin{aligned}
 \mathbb{P}(X \geq n) &\geq 1 - \frac{n^2 p^2 (1 - p)}{(np - 2)^2 E(X)} \text{ (By part (e))} \\
 &= 1 - \frac{2p(1 - p)}{\left(p - \frac{2}{n}\right)^2 (n(n - 1))} \\
 &\rightarrow 1 - 0 = 1 \text{ as } n \rightarrow \infty \text{ since } p \text{ is constant}
 \end{aligned}$$

Therefore, $\mathbb{P}(X \geq n) \rightarrow 1$ as $n \rightarrow \infty$. Also, since a graph with more than $n - 1$ edges must contain a cycle, a.a.s. G contains a cycle.

QUESTION 3

Solution: Consider the graph model $G = G(n, p)$.

(a) Let X be the random variable counting the number of triangles. For all subsets S such that $|S| = 3$ of $V(G)$ define the indicator variable X_S as follows:

$$X_S = \begin{cases} 1 & \text{if } G[S] \cong K_3 \\ 0 & \text{otherwise} \end{cases}$$

Now, $E(X_S) = \mathbb{P}(X_S = 1) = p^3$, so by the linearity of expectation,

$$\begin{aligned} E(X) &= \sum_{S \subseteq V(G), |S|=3} E(X_S) \\ &= \sum_{S \subseteq V(G), |S|=3} p^3 \\ &= \binom{n}{3} p^3 \end{aligned}$$

(b) $Var(X) = Var(\sum X_S) = \sum_{s \subseteq V(G): |s|=3} Var(X_S) + \sum_{S \neq T} Cov(X_S, X_T)$ (*). Now $Var(X_S) = E(X_S^2) - (E(X_S))^2 = p^3 - p^6 \leq p^3$. Also, for $S \neq T$,

$$\begin{aligned} Cov(X_S, X_T) &= E(X_S X_T) - E(X_S)E(X_T) \\ &= \mathbb{P}(X_S X_T = 1) - p^6 \\ &\leq p^5 - p^6 \quad (\text{since 2 triangles may share at most 1 edge}) \\ &\leq p^5 \end{aligned}$$

So from (*),

$$\begin{aligned}
\text{Var}(X) &\leq \sum_{s \subseteq V(G): |S|=3} p^3 + \sum_{S \neq T} p^5 \\
&\leq \binom{n}{3} p^3 + 4 \cdot \binom{n}{4} p^5 \quad (\text{since every subset of 4 vertices contains 4 possible triangles}) \\
&\leq n^3 p^3 + n^4 p^4 \quad (\text{since } p \leq 1)
\end{aligned}$$

(c) Claim: The threshold function for G has a triangle is $f(n) = \frac{1}{n}$.

i) If $\frac{1}{n} \ll p$, then by Chebyshev's inequality:

$$\begin{aligned}
\mathbb{P}(X = 0) &\leq \mathbb{P}(|X - E(X)| \geq E(X)) \\
&\leq \frac{\text{Var}(X)}{E(X)^2} \\
&\leq \frac{n^3 p^3 + n^4 p^4}{\left(\binom{n}{3} p^3\right)^2} \\
&= \theta\left(\frac{1}{n^3 p^3} + \frac{1}{n^2 p^2}\right) \\
&\rightarrow 0 \text{ as } n \rightarrow \infty \text{ (since } \frac{1}{n} \ll p)
\end{aligned}$$

ii) If $p \ll \frac{1}{n}$, then by Markov's inequality:

$$\begin{aligned}
\mathbb{P}(X \geq 1) &\leq E(X) \\
&\leq (np)^3 \\
&\rightarrow 0 \text{ as } n \rightarrow \infty \text{ (since } p \ll \frac{1}{n})
\end{aligned}$$

Therefore, $f(n) = \frac{1}{n}$ is a threshold function for "G has a triangle".

QUESTION 5

For every graph H there exists a function $p = p(n)$ so that $\lim_{n \rightarrow \infty} p(n) = 0$ but a.a.s a graph G produced by $G(n, p)$ contains an induced copy of H .

Proof. We prove the result by the probabilistic method. Fix H , a finite graph. Let X be the number of induced subgraphs of $G(n, p)$ of size $|H|$ that are isomorphic to H . We define the indicator variable X_S for all $S \subseteq V(G)$ with $|S| = |H|$ as follows:

$$X_S = \begin{cases} 1 & \text{if } G[S] \cong H \\ 0 & \text{otherwise} \end{cases}$$

Let $m = \max\{e(H), \binom{|H|}{2} - e(H)\}$. We claim that if $p(n) = \left(\frac{1}{n}\right)^{\frac{|H|-1}{m}}$, then a.a.s G contains an induced copy of H . Now $X = \sum_{S \subseteq V(G): |S|=|H|} X_S$ and $E(X_S) = P(X_S = 1) \geq p^{e(H)}(1-p)^{\binom{|H|}{2}-e(H)}$, so

$$\begin{aligned} E(X) &= E\left(\sum_{S \subseteq V(G): |S|=|H|} X_S\right) \\ &= \sum_{S \subseteq V(G): |S|=|H|} E(X_S) \\ &\geq \sum_{S \subseteq V(G): |S|=|H|} p^{e(H)}(1-p)^{\binom{|H|}{2}-e(H)} \\ &= \binom{n}{|H|} p^{e(H)}(1-p)^{\binom{|H|}{2}-e(H)} \\ &\geq \left(\frac{n}{|H|}\right)^{|H|} p^m(1-p)^m \quad \text{since } p, (1-p) < 1 \\ &= \left(\frac{n}{|H|}\right)^{|H|} \left(\frac{1}{n}\right)^{\frac{m(|H|-1)}{m}} \left(1 - \left(\frac{1}{n}\right)^{\frac{|H|-1}{m}}\right)^m \\ &= \left(\frac{n}{|H|^{|H|}}\right) \left(1 - \left(\frac{1}{n}\right)^{\frac{|H|-1}{m}}\right)^m \end{aligned}$$

And since $|H|$ and m are constants, $\left(1 - \left(\frac{1}{n}\right)^{\frac{|H|-1}{m}}\right)^m \rightarrow 1$ as $n \rightarrow \infty$, so,

$$\left(\frac{n}{|H|^{|H|}}\right) \left(1 - \left(\frac{1}{n}\right)^{\frac{|H|-1}{m}}\right)^m \rightarrow \infty \quad \text{as } n \rightarrow \infty.$$

Therefore, $E(X) \rightarrow \infty$ as $n \rightarrow \infty$.

We now compute the variance of X so that we may apply Chebyshev's Inequality to show that $\mathbb{P}(X = 0) \rightarrow 0$ as $n \rightarrow \infty$.

$$\begin{aligned}
\text{Var}(X) &= \sum_{S \subseteq V(G): |S|=|H|} \text{Var}(X_S) + \sum_{S \neq T} \text{Cov}(X_S, X_T) \\
&= \sum_{S \subseteq V(G): |S|=|H|} \left(E(X_S^2) - E(X_S)^2 \right) + \sum_{S \neq T} \left(E(X_S X_T) - E(X_S)E(X_T) \right) \\
&\leq E(X) - \left(\frac{n}{|H|^{|H|}} \right) \left(1 - \left(\frac{1}{n} \right)^{\frac{|H|-1}{m}} \right)^{2m} + \sum_{S \neq T} \left(E(X_S X_T) \right) - \left(\binom{n}{2} \right) \left(1 - \left(\frac{1}{n} \right)^{\frac{|H|-1}{m}} \right)^{2m} \\
&\leq 2E(X) - \left(\frac{n}{|H|^{|H|}} \right) \left(1 - \left(\frac{1}{n} \right)^{\frac{|H|-1}{m}} \right)^{2m} \quad (*)
\end{aligned}$$

Now, by Chebyshev,

$$\begin{aligned}
\mathbb{P}(X = 0) &\leq \mathbb{P}(|E(X) - X| \geq E(X)) \\
&\leq \frac{\text{Var}(X)}{E(X)^2} \\
&\leq \frac{2E(X) - \left(\frac{n}{|H|^{|H|}} \right) \left(1 - \left(\frac{1}{n} \right)^{\frac{|H|-1}{m}} \right)^{2m}}{E(X)^2} \quad (\text{by } (*)) \\
&\leq \frac{2}{E(X)} - \frac{\left(\frac{n}{|H|^{|H|}} \right) \left(1 - \left(\frac{1}{n} \right)^{\frac{|H|-1}{m}} \right)^{2m}}{E(X)^2} \\
&\rightarrow 0 - 0 = 0 \quad \text{as } n \rightarrow \infty \text{ (since } E(X) \rightarrow \infty \text{ and } \left(\frac{n}{|H|^{|H|}} \right) \left(1 - \left(\frac{1}{n} \right)^{\frac{|H|-1}{m}} \right)^{2m} < E(X))
\end{aligned}$$

Therefore, $\mathbb{P}(X \geq 1) = 1 - \mathbb{P}(X = 0) \rightarrow 1$ as $n \rightarrow \infty$. So a.a.s. G contains a copy of H .

□