

Definition: The **cartesian product** of G and H , written $G \square H$ is the graph with vertex set $V(G) \times V(H)$ specified by putting (u, v) adjacent to (u', v') if and only if

(1) $u = u'$ and $vv' \in E(H)$,

or

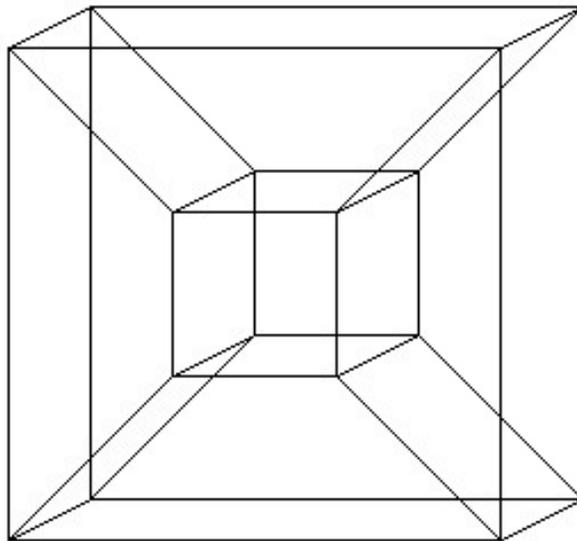
(2) $v = v'$ and $uu' \in E(G)$.

Example: $C_3 \square C_4$

The cartesian product operation is symmetric;

$$G \square H \cong H \square G.$$

The hypercube is another familiar example.



Proposition: $\chi(G \square H) = \max \{ \chi(G), \chi(H) \}$

Proof: The cartesian product $G \square H$ contains copies of G and H as subgraphs, so
 $\chi(G \square H) \geq \max \{ \chi(G), \chi(H) \}$

Let $k = \max \{ \chi(G), \chi(H) \}$. To prove the upper bound, we produce a proper k -colouring of $G \square H$ using optimal colourings of G and H . Let g be a proper $\chi(G)$ -colouring of G , and let h be a proper $\chi(H)$ -colouring of H .

Define a colouring f of $G \square H$ by letting $f(u, v)$ be the congruence class of $g(u) + h(v)$ modulo k . Thus f assigns colours to $V(G \square H)$ from a set of size k .

We claim that f properly colours $G \square H$. If (u, v) and (u', v') are adjacent in $G \square H$, then $g(u) + h(v)$ and $g(u') + h(v')$ agree in one summand and differ by between 1 and k in the other. Since the difference of the two sums is between 1 and k , they lie in different congruence classes modulo k .

The cartesian product allows us to compute chromatic numbers by computing independence numbers, because a graph G is m -colourable if and only if the cartesian product $G \square K_m$ has an independent set of size $n(G)$.