

Agda is based on

- Type Theory
- Lambda calculus
- Curry-Howard isomorphism, BHK interpretation
- Inductive Types

What are the natural numbers?

Definition The \mathbb{N} of natural numbers has the following structure and properties:

- There is an element called "0": $0 : \mathbb{N}$
- There is a function called "succ": $\text{succ} : \mathbb{N} \rightarrow \mathbb{N}$
- Universality: The set \mathbb{N} is the "free" one with the above structure. More precisely:

- Given any (other) set X with an element $z \in X$ and a function $s : X \rightarrow X$,

- There exists a unique function

$$f : \mathbb{N} \rightarrow X$$

such that

$$(1) \quad f(0) = z$$

$$(2) \quad \text{for all } n \in \mathbb{N},$$

$$f(\text{succ}(n)) = s(f(n)).$$

To make this definition easier to understand, let's introduce some concepts.

Definition A zero-successor algebra is a triple (X, z, s) such that

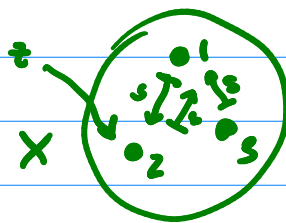
- X is a set
- $z \in X$ is an element called the "zero" of the algebra,
- $s: X \rightarrow X$ is a function called the "successor function" of the algebra.

Examples (1) The set of natural numbers $(\mathbb{N}, 0, \text{succ})$ where $\text{succ}(n) = n+1$ is a zero-successor algebra.

(2) $(\mathbb{N}, 0, f)$ where $f(n) = n+2$ is a zero-successor algebra.

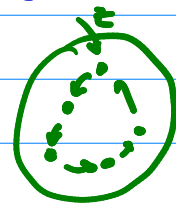
(3) Let $X = \{1, 2, 3\}$. Define $s: X \rightarrow X$ by $s(1) = 2$, $s(2) = 1$, $s(3) = 1$

Then $(X, 2, s)$ is a zero-successor algebra.



(4) Let $X = \{0, 1, 2, 3, 4\}$ Define $z = 0$

and $s(0) = 1$ $s(3) = 4$
 $s(1) = 2$ $s(4) = 0$
 $s(2) = 3$

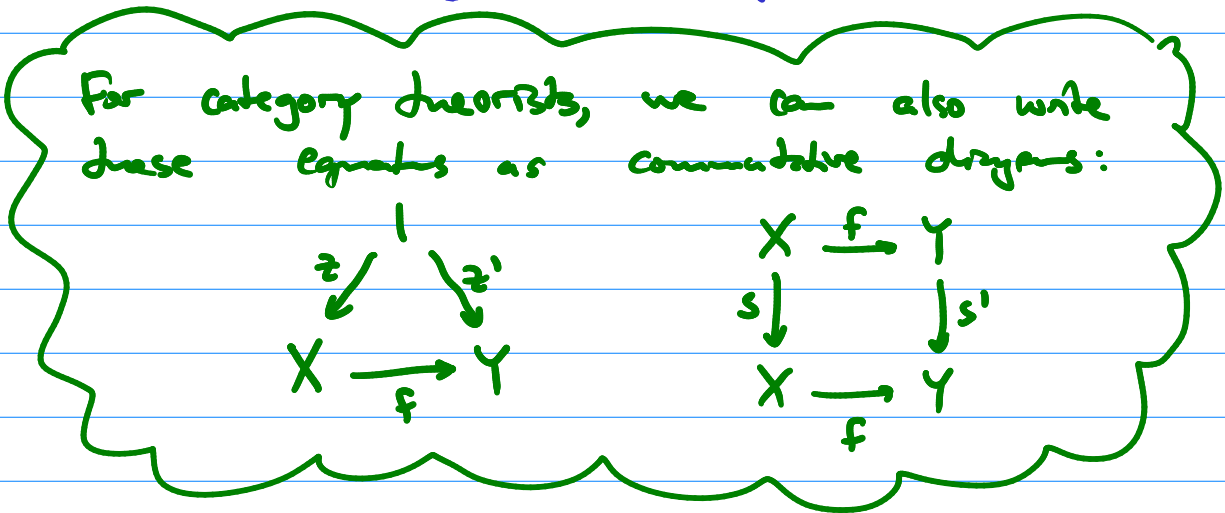


Definition let (X, z, s) and (Y, z', s') be zero-successor algebras. A homomorphism (or zero-successor algebra) is a function

$$f: X \rightarrow Y$$

such that $f(z) = z'$

and $f(s(x)) = s'(f(x))$



Example Consider $(\mathbb{N}, 0, \text{succ})$ and (X, z, s)

where $X = \{a, b, c\}$, $z = a$, $s(a) = b$, $s(b) = c$, $s(c) = a$.

$$f(0) = a$$

$$f(1) = f(\text{succ}(0))$$

$$= s(f(0))$$

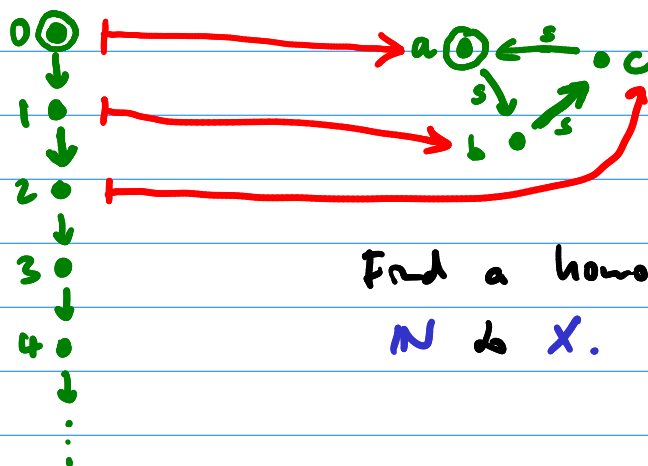
$$= s(a)$$

$$= b$$

$$f(2) = f(\text{succ}(1))$$

$$= s(f(1))$$

$$= s(b) = c$$



Find a homomorphism from \mathbb{N} to X .

Theorem There exists a unique homomorphism of τ -successor algebras $f: \mathbb{N} \rightarrow X$

Proof By induction. $f(0)$ is uniquely determined. And whenever we know $f(n)$, then $f(n+1)$ is uniquely determined. \square

Theorem For any τ -successor algebra (X, z, s) , there exists a unique homomorphism $f: \mathbb{N} \rightarrow X$.

Proof The same. \square

Definition A τ S-algebra (X, z, s) is called initial if it has the property that for all τ S-algebras (Y, z', s') , there exists a unique homomorphism $f: X \rightarrow Y$

Remark $(\mathbb{N}, 0, \text{succ})$ is an initial τ S-algebra.

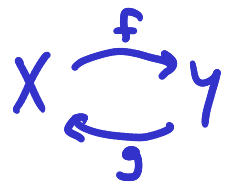
Theorem Any two initial τ S-algebras are isomorphic.

Proof Let (X, z, s) and (Y, z', s') be τ S-algebras, and assume both X and Y are initial.

(a) Since X is initial, there exists a unique homomorphism $f: X \rightarrow Y$.

(b) Since Y is initial, there exists a unique homomorphism $g: Y \rightarrow X$.

(c) Then $g \circ f : X \rightarrow X$
is a homomorphism.



But $id : X \rightarrow X$

is also a homomorphism.

But X is initial. Therefore $g \circ f = id$

(d) Similarly, $f \circ g = id$

(e) So f, g are inverses, so
isomorphisms. So (X, z, s) and (Y, z', s')
are isomorphic. \square

Theorem $(\mathbb{N}, 0, succ)$ are (up to isomorphism)

the unique initial zero-successor algebra.

\square

In Agda when we write down a definition
of a data type

data \mathbb{N} : Set where

zero : \mathbb{N}

succ : $\mathbb{N} \rightarrow \mathbb{N}$

it is automatically understood that \mathbb{N} is the
initial such algebra. This is called an
inductive datatype

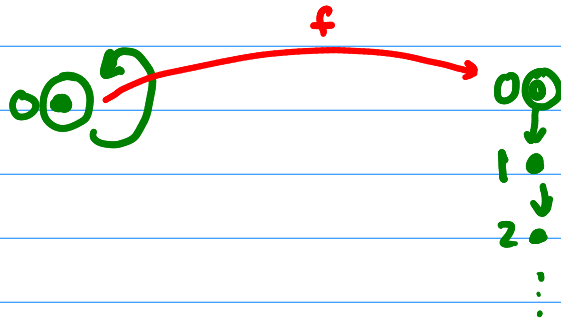
Question: Is (X, z, s) initial, where

$$X = \{0\}, z = 0, s(0) = 0?$$



No, because there exists no homomorphism into \mathbb{N} :

$$(X, z, s) \xrightarrow{f} (\mathbb{N}, 0, \text{succ}).$$



$$f(0) = 0$$

$$f(s(0)) = 0$$

$$\text{succ}(f(0)) = 1$$

Not a homo.