

A quick introduction of constructive logic

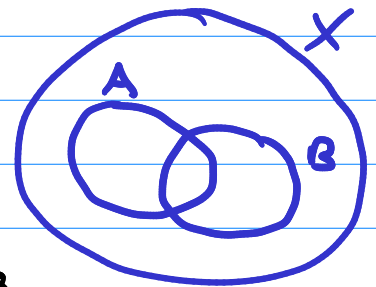
I. Classical logic

1. - Truth tables: Every propositional variable is interpreted as "T" or "F".
 - Value of compound formulas is given by truth tables.

A	B	$A \vee B$
T	T	T
T	F	T
F	T	T
F	F	F

2. - Venn diagrams Given a set X called the "universe".

- Every propositional variable is assigned a subset of X .



- The logical connectives are interpreted as operations on sets:

$$A \text{ and } B = A \cap B \quad \text{"intersection"}$$

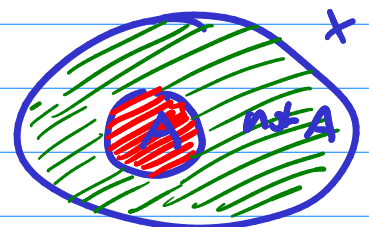
$$A \text{ or } B = A \cup B \quad \text{"union"}$$

$$\text{not } A = X - A = A^c \quad \text{"complement"}$$

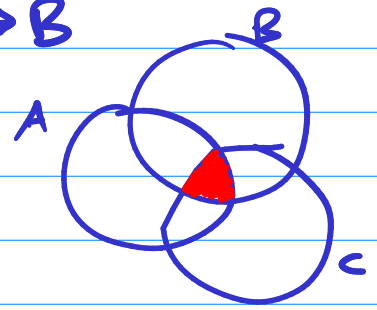
- A proposition P is a tautology or valid if $P = X$.

ex. $A \vee \neg A$

$$A \cup A^c = X.$$



$$(A \cap B) \cap C \rightarrow B$$

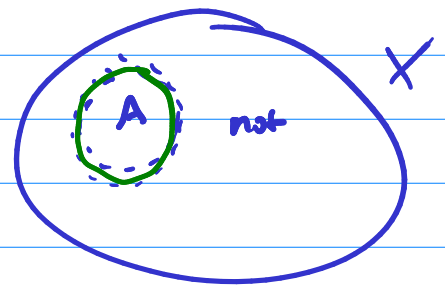


is valid because

$$(A \cap B) \cap C \subseteq B.$$

II. Constructive logic

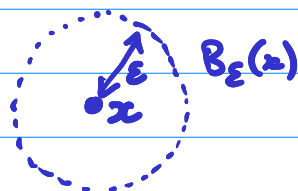
Same idea, but use open sets.



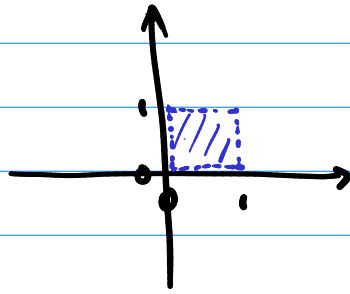
A model of intuitionistic (propositional) logic
[Stone 1937, Tarski 1938]

Definition A subset A of \mathbb{R}^2 is called open if for every $x \in A$, there exists $\epsilon > 0$ such that $B_\epsilon(x) \subseteq A$.

Here $B_\epsilon(x) = \{y \in \mathbb{R}^2 \mid \|x - y\| < \epsilon\}$
is the ϵ -neighborhood of x .



Examples let $A = \{ (x, y) \mid 0 < x < 1 \text{ and } 0 < y < 1 \}$



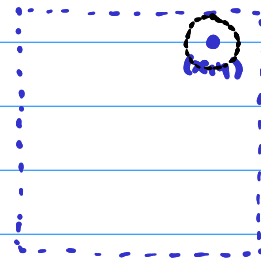
Then A is open.

Proof: Consider any $(x, y) \in A$.

Then $0 < x < 1$ and $0 < y < 1$.

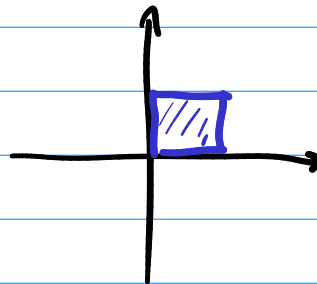
let $\epsilon = \min \{ x, 1-x, y, 1-y \}$

Then $B_\epsilon(x, y) \subseteq A$

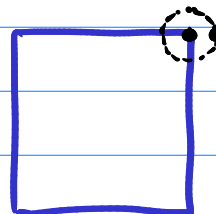


let $B = \{ (x, y) \mid 0 \leq x \leq 1, 0 \leq y \leq 1 \}$

Then B is not open.



Proof: E.g. $(1, 1) \in B$, but there is no $\epsilon > 0$ such that $B_\epsilon(1, 1) \subseteq B$.



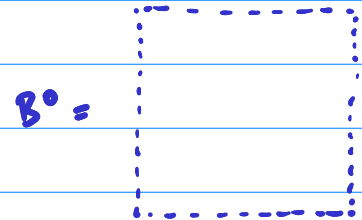
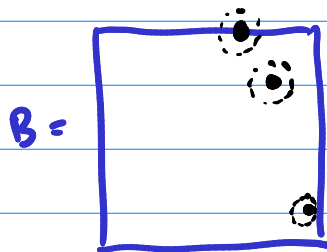
Definition Let $A \subseteq \mathbb{R}^2$ be any subset of \mathbb{R}^2 .
 The interior of A , in symbols A° , is
 defined as follows:

$$A^\circ = \{x \in A \mid \exists \varepsilon > 0. B_\varepsilon(x) \subseteq A\}.$$

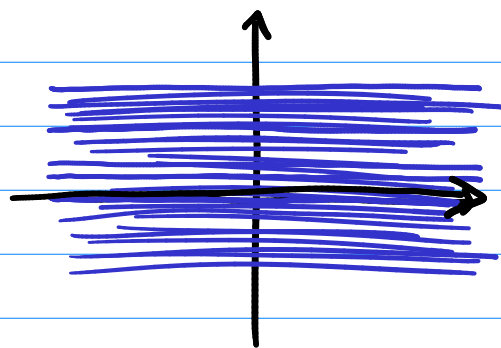
Equivalently, one point x is in the interior of A
 iff A contains some ε -neighborhood of x .

Example let $B = \{(x, y) \mid 0 \leq x \leq 1, 0 \leq y \leq 1\}$.

Then $B^\circ = \{(x, y) \mid 0 < x < 1, 0 < y < 1\}$.



Example let $C = \{(x, y) \mid x \in \mathbb{R}, y \in \mathbb{Q}\}$



Is C open?

No. For example
 $(0, 0) \in C$.

But for all $\varepsilon > 0$,
 there exists some
 irrational number
 $y \in (0, \varepsilon)$.

Then $(0, y) \in B_\varepsilon(0, 0)$
 but $(0, y) \notin C$.

The interior of C is $C^\circ = \emptyset$.

Proposition If $A \subseteq \mathbb{R}^2$ is any set, then A° is open.

Proof: Exercise.

The Stone-Tarski model of intuitionistic logic

Consider the language with propositional symbols

A_1, A_2, A_3, \dots

and with the usual logical connectives

$\wedge, \vee, \neg, \rightarrow, \top, \perp$

So in other words, the propositions of the logic are given by the grammar

$$A, B ::= A_i \mid (A \wedge B) \mid (A \vee B) \mid \neg A \mid (A \rightarrow B) \mid \top \mid \perp$$

Definition An interpretation of the logic is a function σ that assigns to each A_i an open subset $\sigma(A_i) \subseteq \mathbb{R}^2$.

Moreover, we extend the interpretation to all propositions as follows:

$$\bar{\sigma}(A_i) = \sigma(A_i)$$

$$\bar{\sigma}(A \wedge B) = \bar{\sigma}(A) \cap \bar{\sigma}(B)$$

$$\bar{\sigma}(A \vee B) = \bar{\sigma}(A) \cup \bar{\sigma}(B)$$

$$\bar{\varepsilon}(T) = \mathbb{R}^2$$

$$\bar{\varepsilon}(F) = \emptyset$$

$$\bar{\varepsilon}(\neg A) = ((\bar{\varepsilon}(A))^c)^{\circ}$$

This is the interior of the complement of $\bar{\varepsilon}(A)$. Note that the complement of an open set may not be open.

$$\bar{\varepsilon}(A \rightarrow B) = (\bar{\varepsilon}(A)^c \cup \bar{\varepsilon}(B))^{\circ}$$

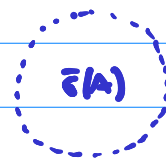
Note: The intersection and the union of two open sets are open. So all the sets defined above are open.

Proposition If a proposition A is intuitionistically valid, then for all interpretations ε , we have $\bar{\varepsilon}(A) = \mathbb{R}^2$.

The converse is also true, i.e. if A is not intuitionistically valid, then there exists some interpretation such that $\bar{\varepsilon}(A) \neq \mathbb{R}^2$.

Examples

1. $A \rightarrow A$



$$\begin{aligned}\bar{\varepsilon}(A \rightarrow A) &= (\bar{\varepsilon}(A)^c \cup \bar{\varepsilon}(A))^{\circ} \\ &= (\mathbb{R}^2)^{\circ} \\ &= \mathbb{R}^2.\end{aligned}$$

$$2. \quad A \vee \neg A$$

$$\overline{\overline{A}} = X = B_1(0,0) \\ = \{(x,y) \mid x^2 + y^2 < 1\}$$

$$\overline{(A \vee \neg A)} = \overline{A} \cup \overline{\neg A} \\ = \overline{A} \cup (\overline{A})^c \\ = X \cup (X^c)^o$$

$$X = \{(x,y) \mid x^2 + y^2 < 1\}$$

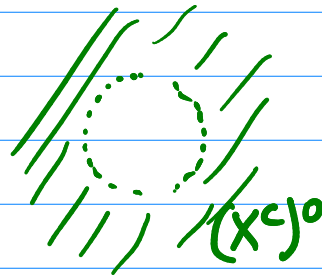
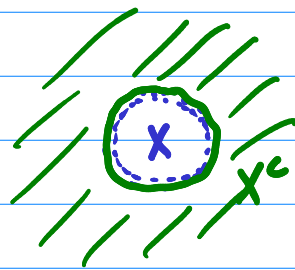
$$X^c = \{(x,y) \mid x^2 + y^2 \geq 1\}$$

$$(X^c)^o = \{(x,y) \mid x^2 + y^2 > 1\}$$

$$X \cup (X^c)^o = \{(x,y) \mid x^2 + y^2 \neq 1\}$$

$$= \mathbb{R}^2 - \{(x,y) \mid x^2 + y^2 = 1\}$$

$$\neq \mathbb{R}^2.$$



$$X \cup (X^c)^o =$$



$$= \mathbb{R}^2 - \text{circle.}$$

$$3. \quad \neg \neg A \text{ vs. } A.$$

Example $\overline{\overline{A}} = X = \{(x,y) \mid x^2 + y^2 < 1\}$

$$X^c = \{(x,y) \mid x^2 + y^2 \geq 1\}$$

$$\overline{\neg A} = (X^c)^o = \{(x,y) \mid x^2 + y^2 > 1\}$$

$$((X^c)^o)^c = \{(x,y) \mid x^2 + y^2 \leq 1\}$$

$$\overline{\neg \neg A} = (((X^c)^o)^c)^o = \{(x,y) \mid x^2 + y^2 < 1\}$$

Exemple 2 $\bar{A} = X = \{ (x, y) \mid x \in \mathbb{R}, y \neq 0 \}$

$$X^c = \{ (x, y) \mid x \in \mathbb{R}, y = 0 \}$$

$$\bar{(\neg A)} = (X^c)^0 = \emptyset$$

$$((X^c)^0)^c = \mathbb{R}^2$$

$$\bar{(\neg \neg A)} = (((X^c)^0)^c)^0 = \mathbb{R}^2$$

$$\bar{A} \neq \bar{(\neg \neg A)}.$$